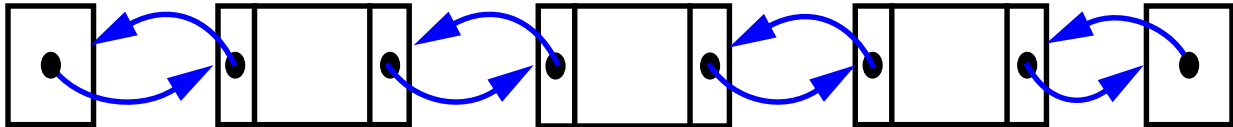


SEQUENCES

- Vectors
- Positions
- Lists
- General Sequences
- Bubble Sort Algorithm



The Vector ADT

- A sequence S (with n elements) that supports the following methods:
 - **elemAtRank(r):**
Return the element of S with rank r ; an error occurs if $r < 0$ or $r > n - 1$
 - **replaceAtRank(r, e):**
Replace the element at rank r with e and return the old element; an error condition occurs if $r < 0$ or $r > n - 1$
 - **insertAtRank(r, e):**
Insert a new element into S which will have rank r ; an error occurs if $r < 0$ or $r > n - 1$
 - **removeAtRank(r):**
Remove from S the element at rank r ; an error occurs if $r < 0$ or $r > n - 1$

Array-Based Implementation

- Some Pseudo-Code:

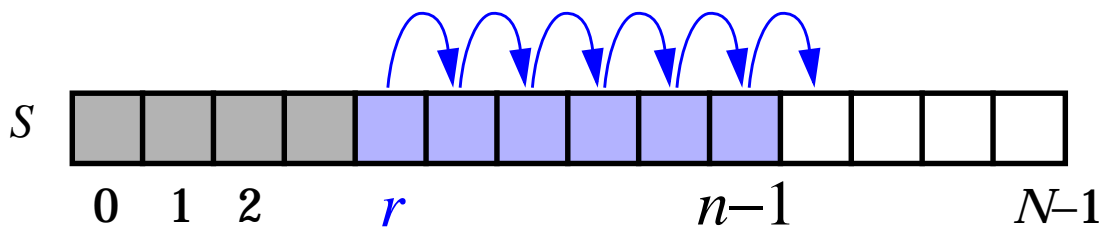
Algorithm insertAtRank(r, e):

for $i = n - 1, n - 2, \dots, r$ **do**

$S[i+1] \leftarrow s[i]$

$S[r] \leftarrow e$

$n \leftarrow n + 1$



Algorithm removeAtRank(r):

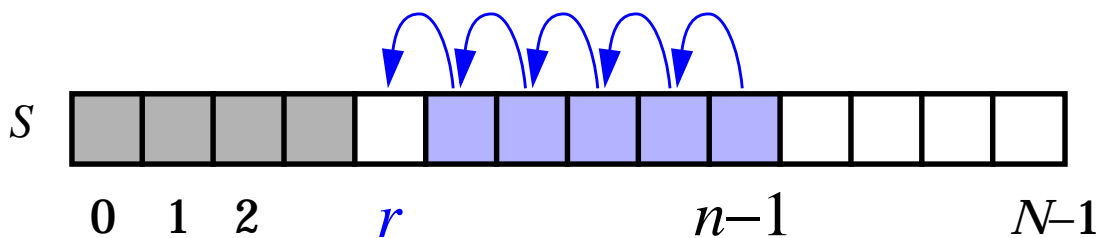
$e \leftarrow S[r]$

for $i = r, r + 1, \dots, n - 2$ **do**

$S[i] \leftarrow S[i + 1]$

$n \leftarrow n - 1$

return



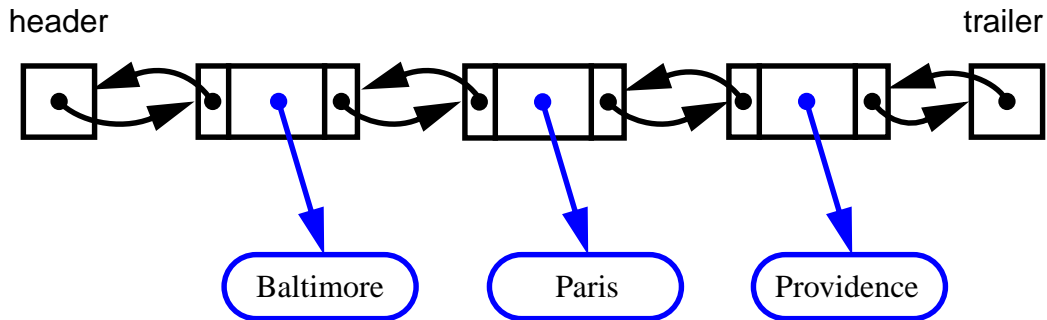
Array-Based Implementation (contd.)

- Time complexity of the various methods:

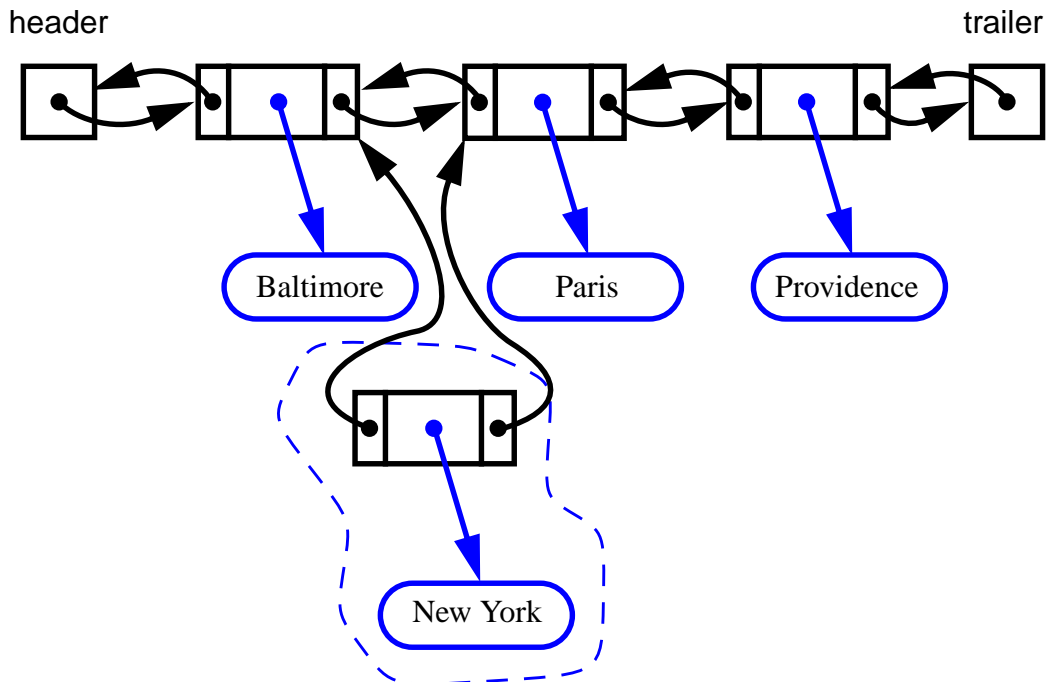
Method	Time
size	$O(1)$
isEmpty	$O(1)$
elemAtRank	$O(1)$
replaceAtRank	$O(1)$
insertAtRank	$O(n)$
removeAtRank	$O(n)$

Implementation with a Doubly Linked List

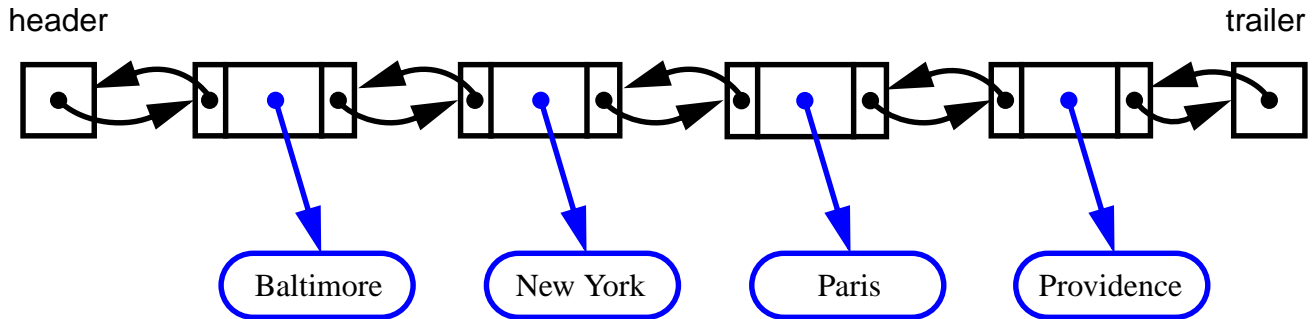
- the list before insertion:



- creating a new node for insertion:



- the list after insertion:



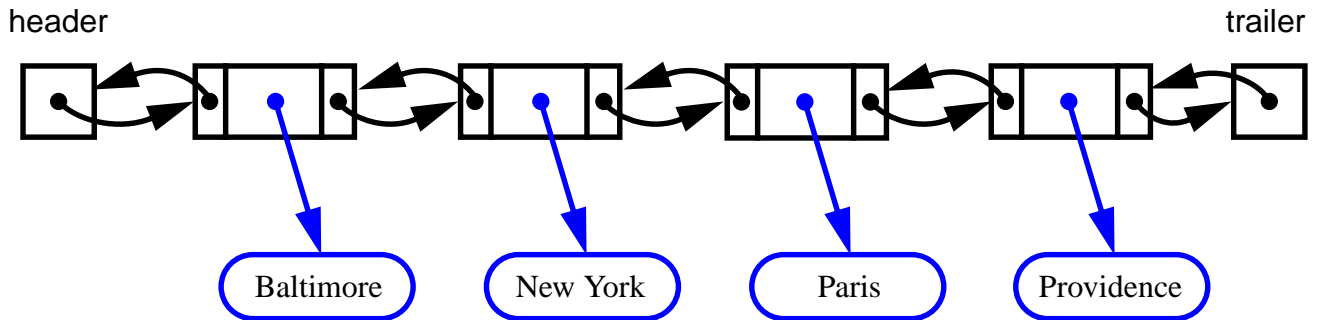
```

public void insertAtRank (int rank, Object element)
    throws BoundaryViolationException {
    if (rank < 0 || rank > size())
        throw new BoundaryViolationException("invalid rank");
    DLNode next = nodeAtRank(rank); // the new node
        //will be right before this
    DLNode prev = next.getPrev(); // the new node
        //will be right after this
    DLNode node = new DLNode(element, prev, next);
    // new node knows about its next & prev. Now
    // we tell next & prev about the new node.
    next.setPrev(node);
    prev.setNext(node);
    size++;
}

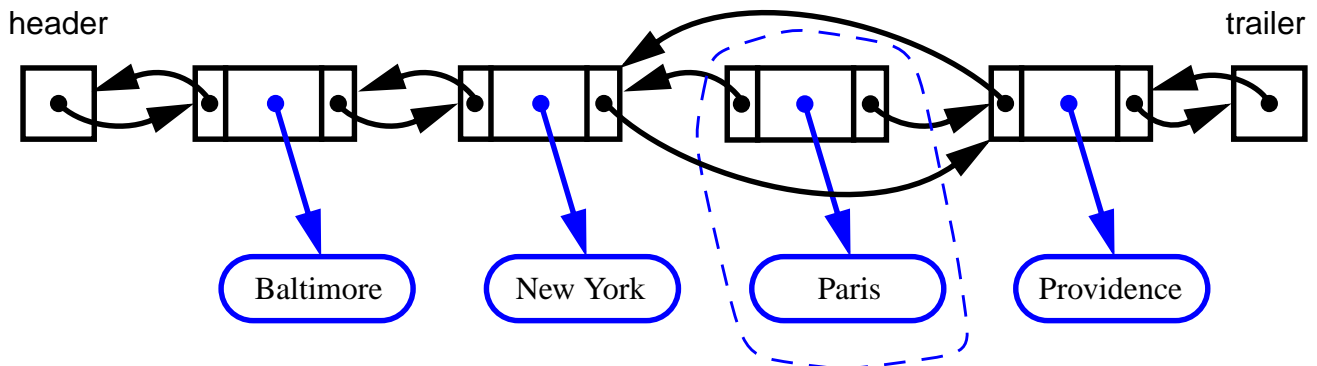
```

Implementation with a Doubly Linked List

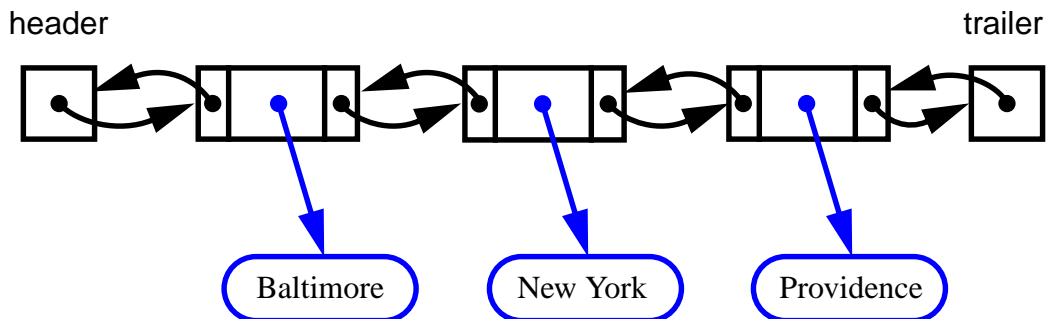
- the list before deletion:



- deleting a node:



- after deletion:



Java Implementation

- code for deletion of a node

```
public Object removeAtRank (int rank)
    throws BoundaryViolationException {
    if (rank < 0 || rank > size()-1)
        throw new BoundaryViolationException("Invalid
            rank.");
    DLNode node = nodeAtRank(rank); // node to
                                    // be removed
    DLNode next = node.getNext(); // node before it
    DLNode prev = node.getPrev(); // node after it
    prev.setNext(next);
    next.setPrev(prev);
    size--;
    return node.getElement(); // returns the
                               // element of the deleted node
}
```


Java Implementation (cont.)

- code for finding a node at a certain rank

```
private DLNode nodeAtRank (int rank) {  
    // auxiliary method to find the node of the  
    // element with the given rank. We make  
    // auxiliary methods private or protected.  
    DLNode node;  
    if (rank <= size()/2) { //scan forward from head  
        node = header.getNext();  
        for (int i=0; i < rank; i++)  
            node = node.getNext();  
    }  
    else { // scan backward from the tail  
        node = trailer.getPrev();  
        for (int i=0; i < size()-rank-1 ; i++)  
            node = node.getPrev();  
    }  
    return node;  
}
```

Nodes

- Linked lists support the efficient execution of *node-based operations*:
 - `removeAtNode(Node v)` and `insertAfterNode(Node v, Object e)`, would be $O(1)$.
- However, node-based operations are not meaningful in an array-based implementation because there are no nodes in an array.
- Nodes are implementation-specific.
- **Dilemma:**
 - If we do not define node based operations, we are not taking full advantage of doubly-linked lists.
 - If we do define them, we violate the generality of ADTs.

From Nodes to Positions

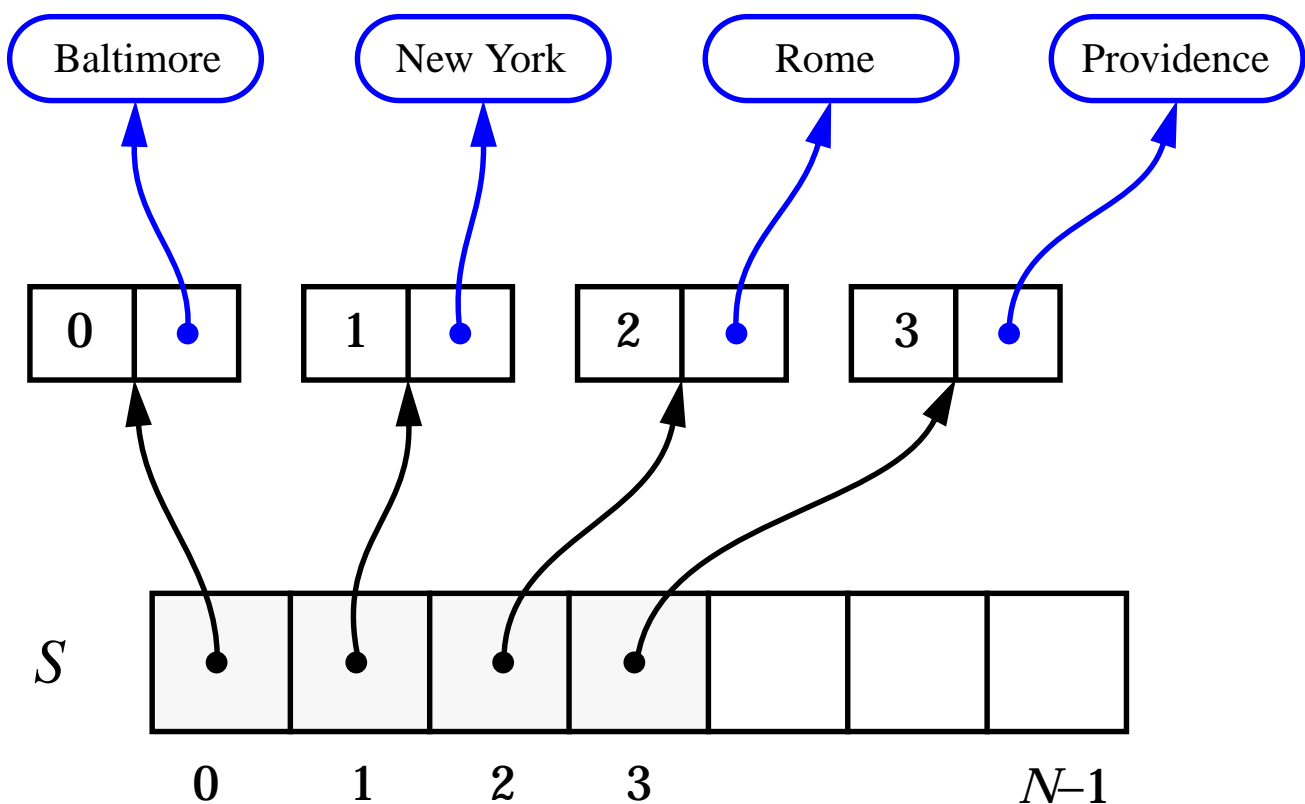
- We introduce the *Position* ADT
- Intuitive notion of “place” of an element.
- Positions have only one method:
`element()`: Return the element at this position
- Positions are defined relatively to other positions (before/after relation)
- Positions are not tied to an element or rank

The List ADT

- ADT with position-based methods
- generic methods `size()`, `isEmpty()`
- query methods `isFirst(p)`, `isLast(p)`
- accessor methods `first()`, `last()`, `before(p)`, `after(p)`
- update methods `swapElements(p,q)`,
`replaceElement(p,e)`, `insertFirst(e)`, `insertLast(e)`,
`insertBefore(p,e)`, `insertAfter(p,e)`, `remove(p)`
- each method takes $O(1)$ time if implemented with a doubly linked list

The Sequence ADT

- Combines the Vector and List ADT (multiple inheritance)
- Adds methods that bridge between ranks and positions
 - `atRank(r)` returns a position
 - `rankOf(p)` returns an integer rank
- An array-based implementation needs to use objects to represent the positions



Comparison of Sequence Implementations

Operations	Array	List
size, isEmpty	$O(1)$	$O(1)$
atRank, rankOf, elemAtRank	$O(1)$	$O(n)$
first, last	$O(1)$	$O(1)$
before, after	$O(1)$	$O(1)$
replaceElement, swapElements	$O(1)$	$O(1)$
replaceAtRank	$O(1)$	$O(n)$
insertAtRank, removeAtRank	$O(n)$	$O(n)$
insertFirst, insertLast	$O(1)$	$O(1)$
insertAfter, insertBefore	$O(n)$	$O(1)$
remove	$O(n)$	$O(1)$

Iterators

- Abstraction of the process of scanning through a collection of elements
- Encapsulates the notions of “place” and “next”
- Extends the position ADT
- Generic and specialized iterators
- *ObjectIterator*
 - hasNext()
 - nextObject()
 - object()
- *PositionIterator*
 - nextPosition()
- Useful methods that return iterators:
 - elements()
 - positions()