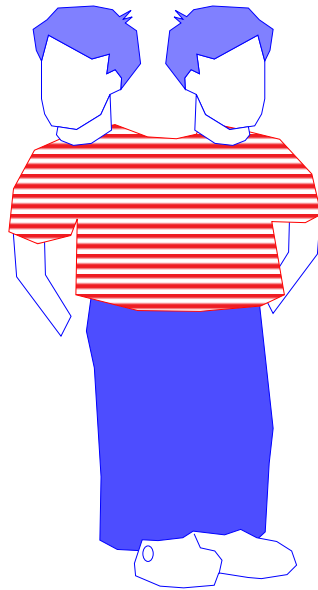


Connectivity and Biconnectivity

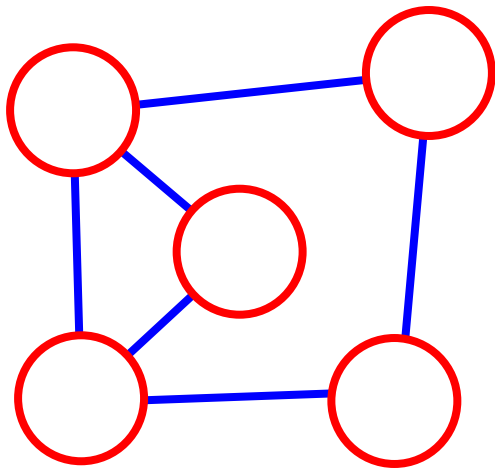


- connected components
- cutvertices
- biconnected components

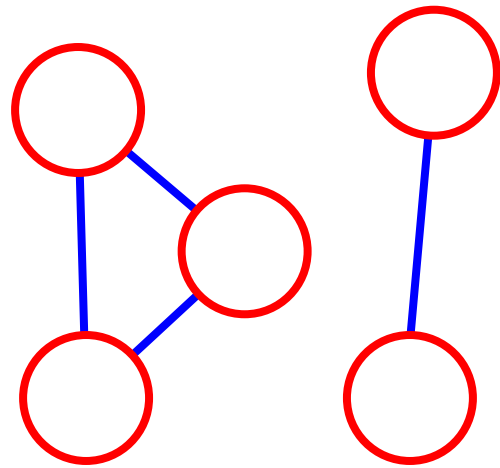


Connected Components

Connected Graph: any two vertices connected by a path



connected



not connected

Connected Component:
maximal connected subgraph of
a graph



Equivalence Relations

A *relation* on a set S is a set R of ordered pairs of elements of S defined by some property

Example:

- $S = \{1,2,3,4\}$
- $R = \{(i,j) \in S \times S \text{ such that } i < j\}$
 $= \{(1,2),(1,3),(1,4),(2,3),(2,4),(3,4)\}$

An *equivalence relation* is a relation with the following properties:

- $(x,x) \in R, \forall x \in S$ (*reflexive*)
- $(x,y) \in R \Rightarrow (y,x) \in R$ (*symmetric*)
- $(x,y), (y,z) \in R \Rightarrow (x,z) \in R$ (*transitive*)

The relation C on the set of vertices of a graph:

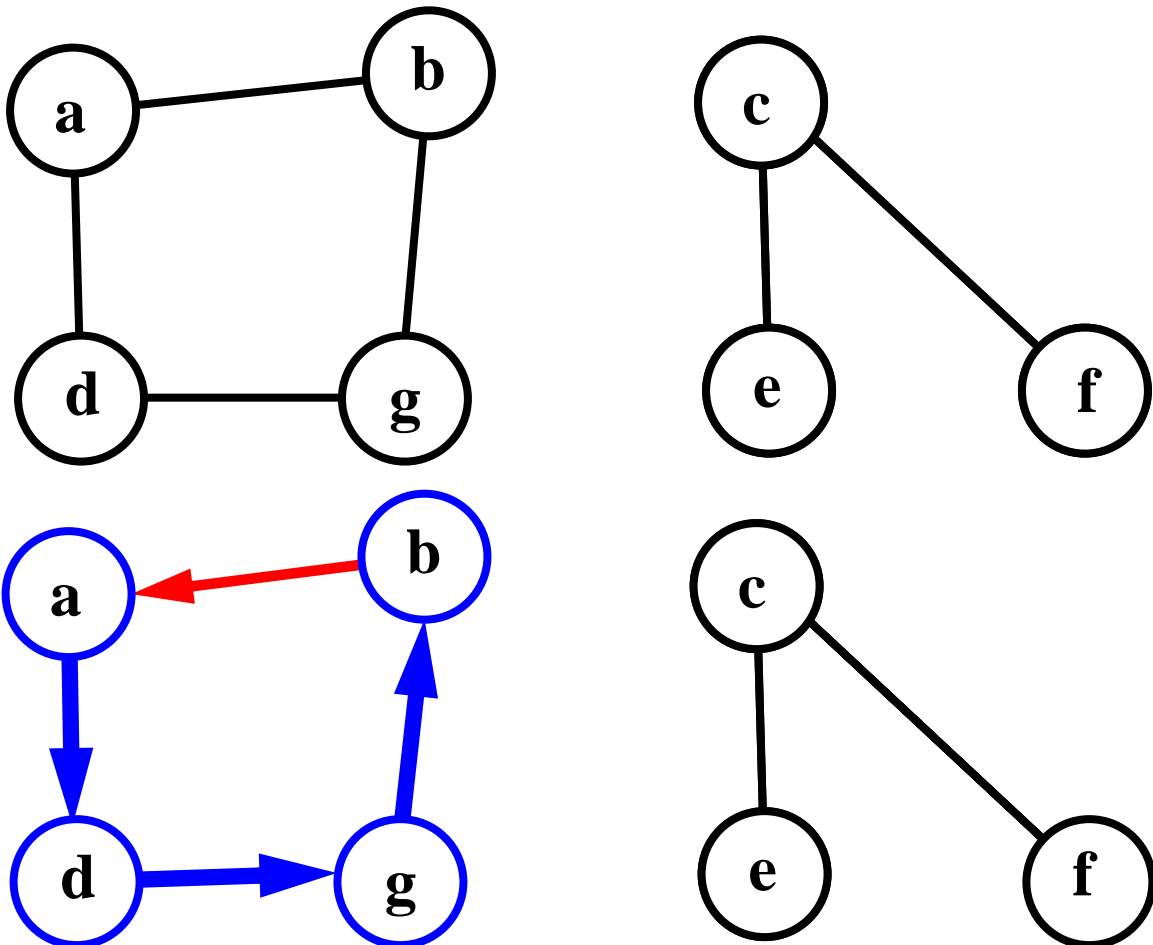
- $(u,v) \in C \Leftrightarrow u$ and v are in the same connected component

is an equivalence relation.



DFS on a Disconnected Graph

- DFS(v) visits all the vertices and edges in the connected component of v



- To compute the connected components:

```

k = 0 // component counter
foreach (vertex v)
  if unvisited(v)
    // add to component k
    // the vertices reached by v
    DFS(v, k++)
  
```

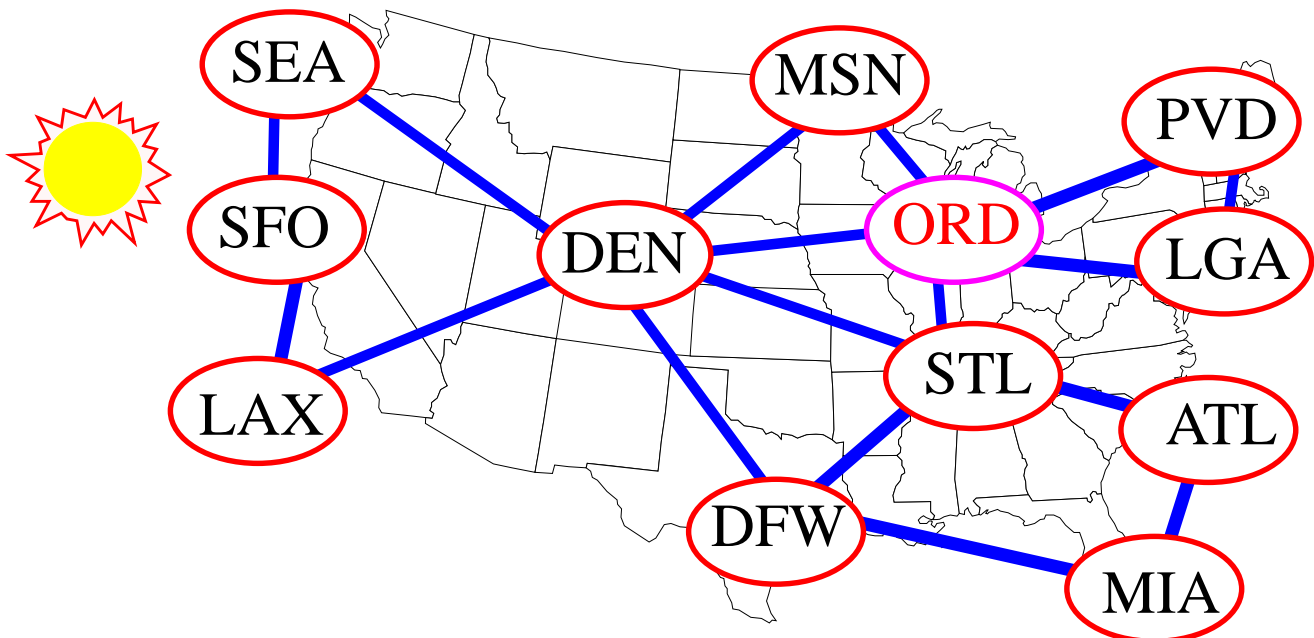


Cutvertices

**Cutvertex (separation vertex):
its removal disconnects the graph**

If the **Chicago** airport is closed, then there is no way to get from Providence to cities on the west coast.

Similarly for **Denver**.

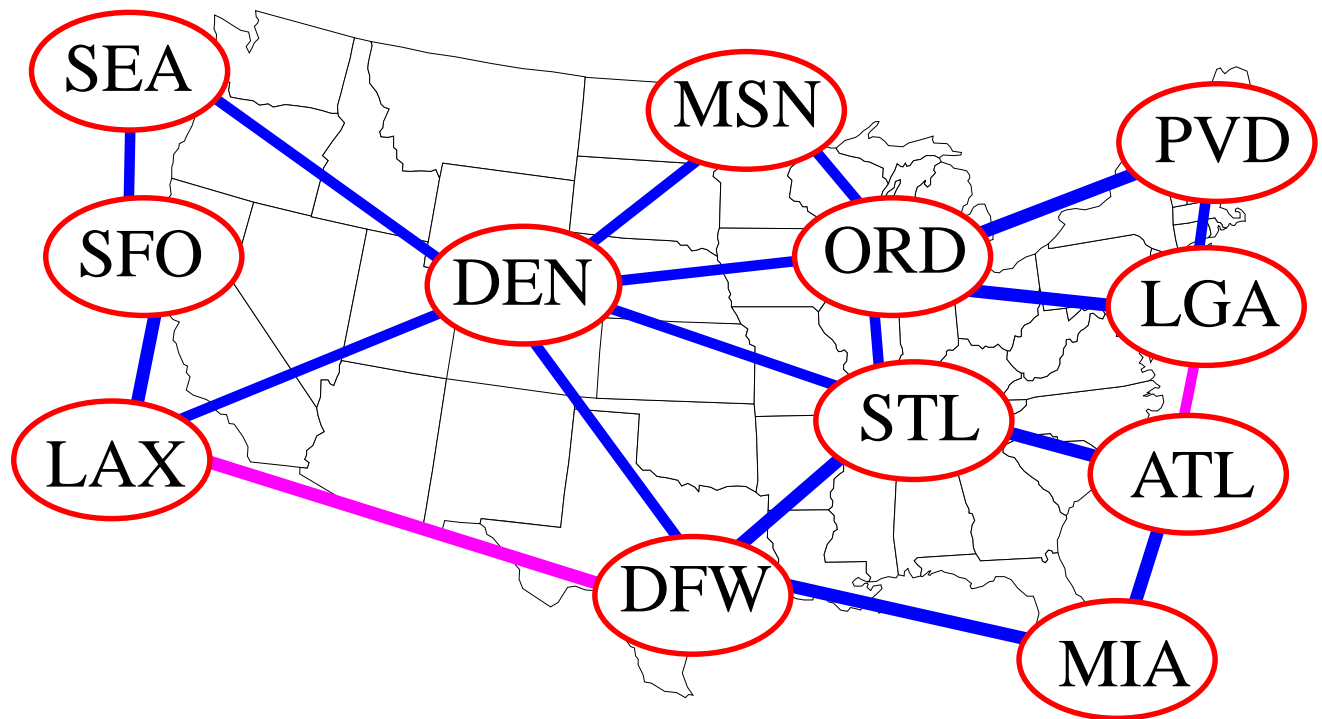


- Cutvertices: **ORD. DEN**



Biconnectivity

Biconnected graph: has no cutvertices

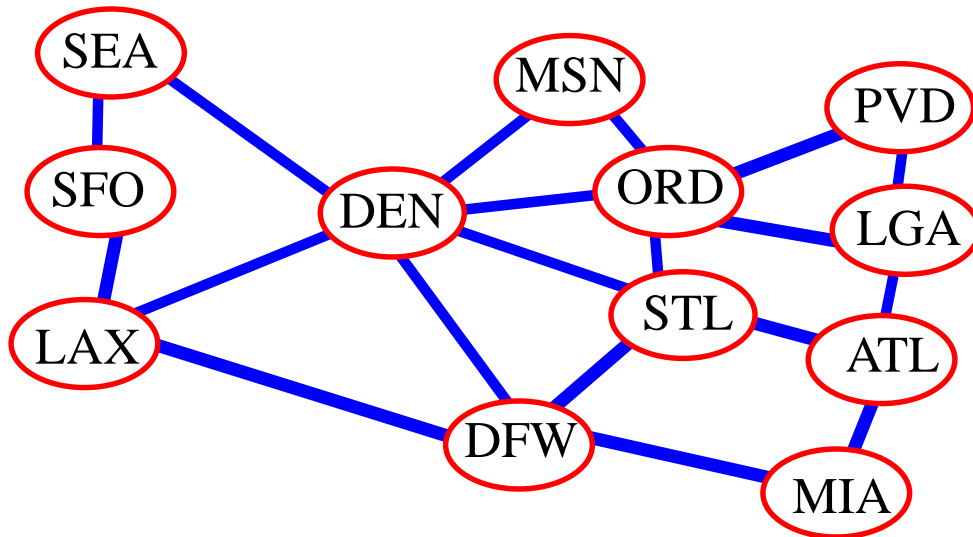


New flights:

LGA-ATL and **DFW-LAX**
make the graph biconnected.



Properties of Biconnected Graphs



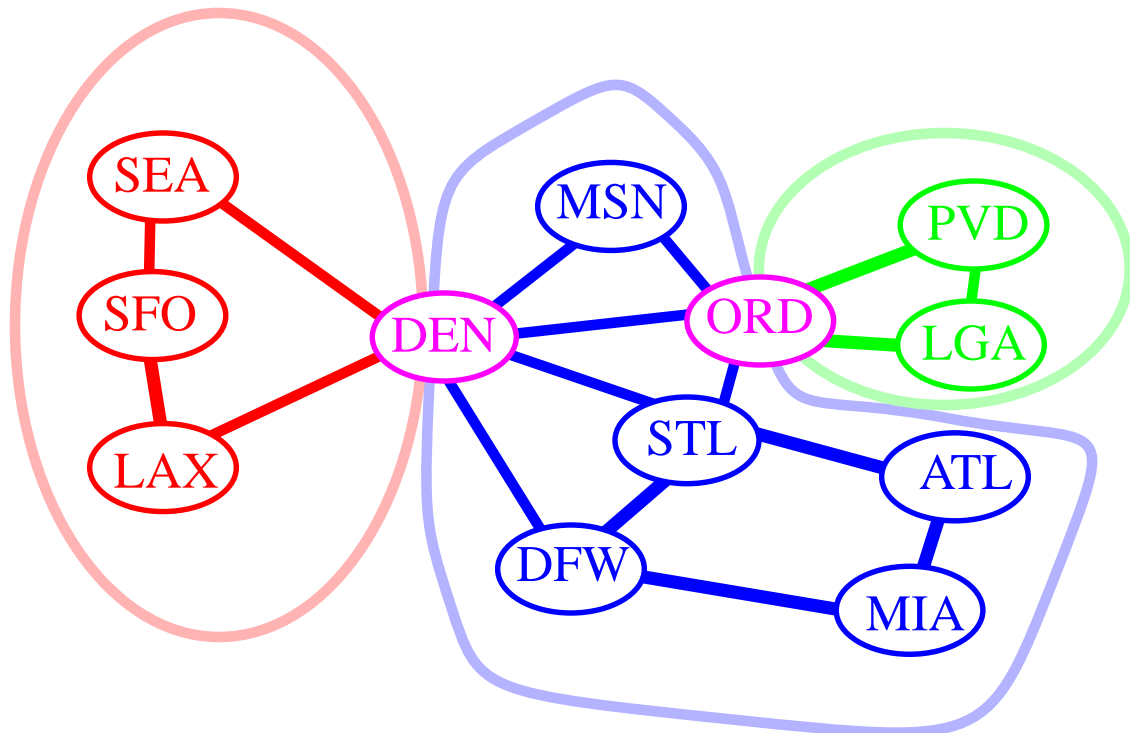
- There are *two disjoint paths* between any two vertices.
- There is a *cycle* through any two vertices.

By convention, two nodes connected by an edge form a biconnected graph, but this does not verify the above properties.



Biconnected Components

- Biconnected component (block):
maximal biconnected subgraph

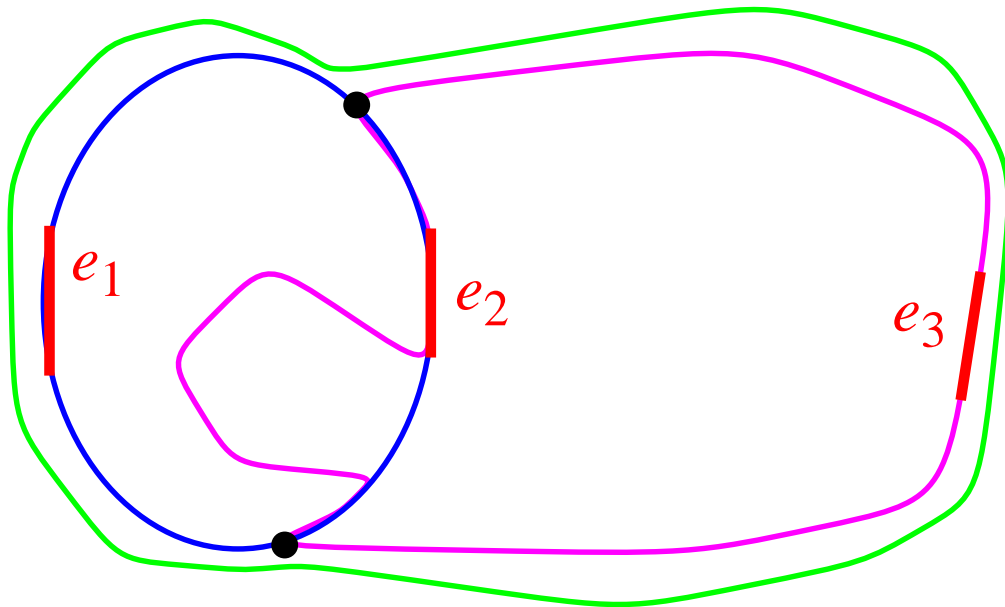


- Biconnected components are edge-disjoint but share **cutvertices**.



Characterization of the Biconnected Components

- *Equivalence relation* R on the *edges* of G :
 $(e', e'') \in R$ if there is a cycle containing both e' and e''
- Proof of the *transitive property*

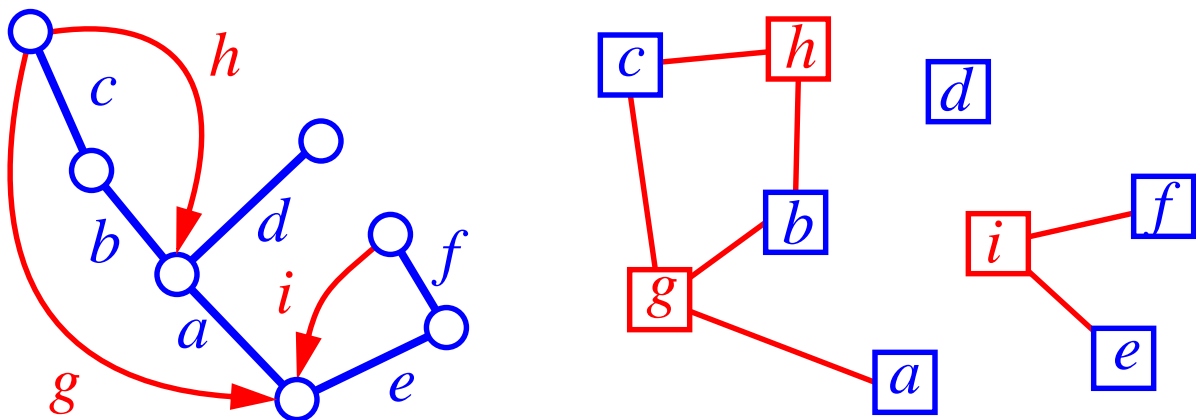


- We partition the edges of G into *equivalence classes* with respect to R .
- Each equivalence class corresponds to
 - a biconnected components of G
 - a connected components of a graph H whose vertices are the *edges* of G and whose edges are the *pairs* in relation R .



DFS and Biconnected Components

- Graph H has $O(m^2)$ edges in the worst case.
- Instead of computing the entire graph H , we use a smaller *proxy* graph K .
- Start with an empty graph K whose vertices are the edges of G .
- Given a DFS on G , consider the $(m - n + 1)$ cycles of G induced by the back edges.
- For each such cycle $C = (e_0, e_1, \dots, e_p)$ add edges $(e_0, e_1) \dots (e_0, e_p)$ to K .

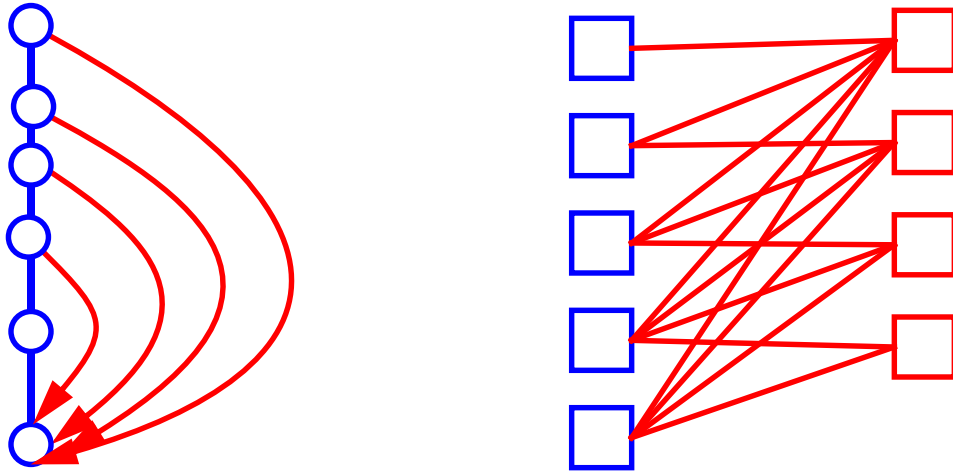


- The connected components of K are the same as those of H !



A Linear Time Algorithm

- The size of K is $O(mn)$ in the worst case.

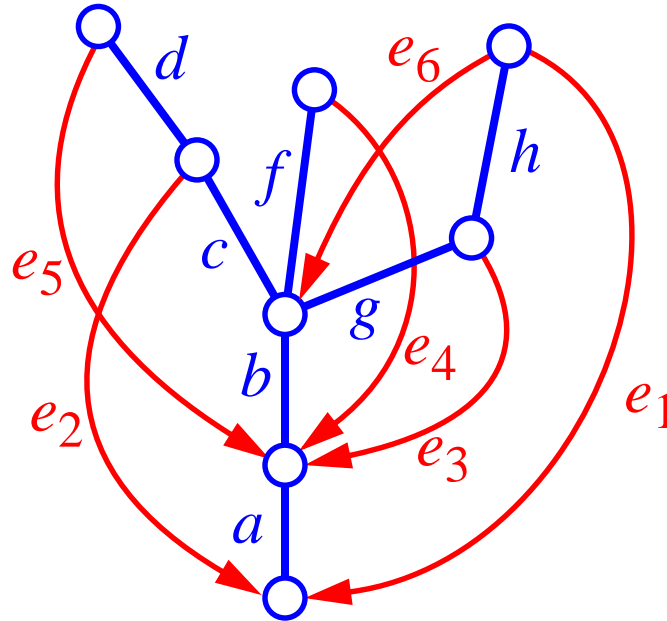


- We can further reduce the size of the proxy graph to $O(m)$
- Process the back edges according to a *preorder visit* of their destination vertex in the DFS tree
- Mark the discovery edges forming the cycles
- Stop adding edges to the proxy graph after the first marked edge is encountered.
- The resulting proxy graph is a forest!
- This algorithm runs in $O(n+m)$ time.

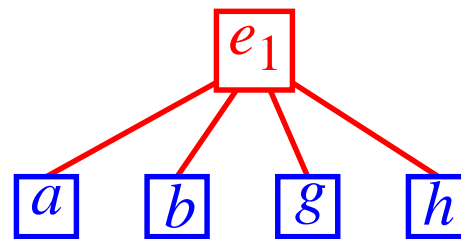


Example

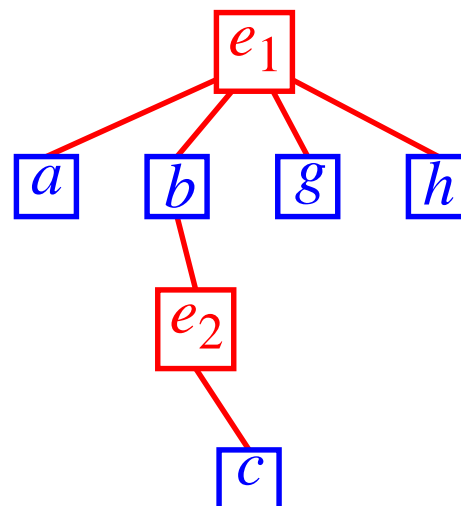
- Back edges labeled according to the preorder visit of their destination vertex in the DFS tree



- Processing e_1

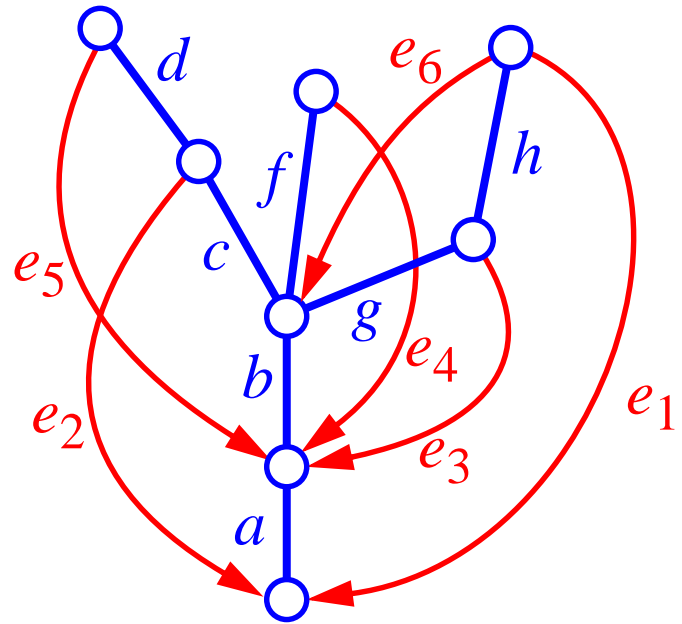


- Processing e_2

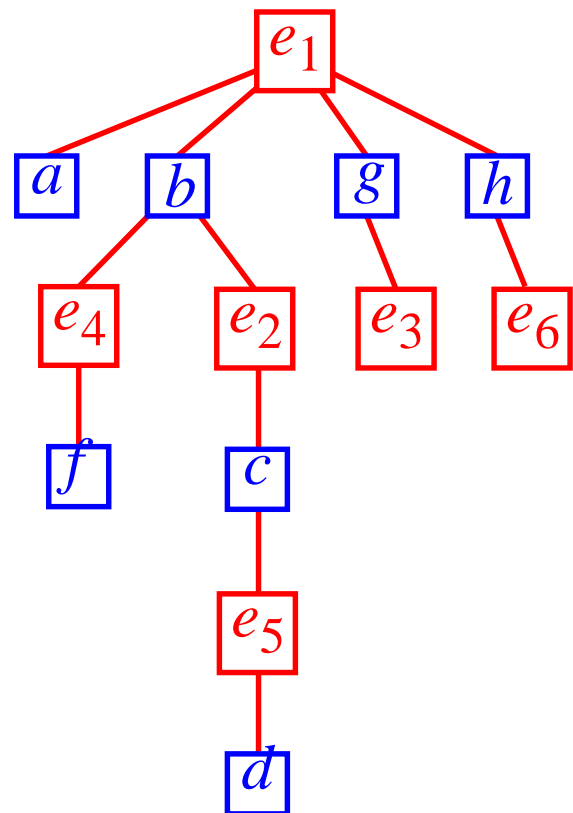


Example (contd.)

- DFS tree

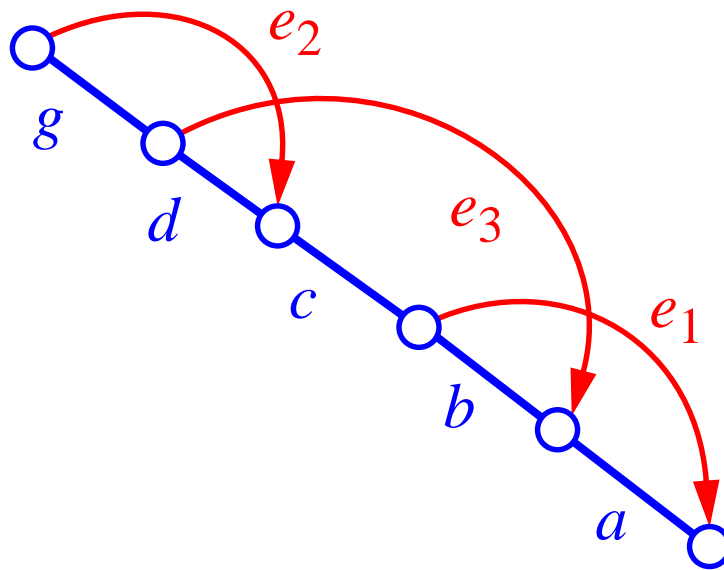


- final proxy graph (a tree since the graph is biconnected)

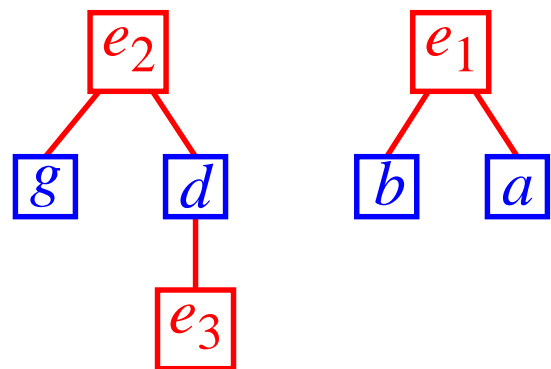


Why Preorder?

- The order in which the back edges are processed is essential for the correctness of the algorithm
- Using a different order ...



- ... yields a graph that provides incorrect information



Try the Algorithm on this Graph!

