

# Symmetric Incoherent Eavesdropping against MDI QKD

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- 1 Review of Quantum Information

# Outline

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- 2 Key Distribution using BB84
  - The Protocol
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## Basic Quantum Algebra for a Single Particle

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- Every allowed reversible physical **transformation** on the state of a particle is represented by a  $2 \times 2$  unitary matrix  $U$ .
- The basis  $\left\{ \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \end{pmatrix} \right\}$  is called the **computational basis** and is represented by  $\{|0\rangle, |1\rangle\}$ .

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## From One to Many Particles

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- For two particles in states  $(\alpha_1|0\rangle + \beta_1|1\rangle)$  and  $(\alpha_2|0\rangle + \beta_2|1\rangle)$ , the joint state is given by the **tensor product**  $(\alpha_1|0\rangle + \beta_1|1\rangle) \otimes (\alpha_2|0\rangle + \beta_2|1\rangle)$   
 $= \alpha_1\alpha_2|00\rangle + \alpha_1\beta_2|01\rangle + \beta_1\alpha_2|10\rangle + \beta_1\beta_2|11\rangle$ ,  
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which is a complex vector in 4 dimensions.
- In general, the state for  $n$  particles is a complex vector in  $2^n$ -dimensional complex vector space (called a **Hilbert Space**).

# Conjugate Bases

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- For  $N = 2$ , the bases  $\{|\psi_1\rangle, |\psi_2\rangle\}$  and  $\{|\phi_1\rangle, |\phi_2\rangle\}$  are conjugate, if and only if we have
$$|\langle\psi_1|\phi_1\rangle|^2 = |\langle\psi_1|\phi_2\rangle|^2 = |\langle\psi_2|\phi_1\rangle|^2 = |\langle\psi_2|\phi_2\rangle|^2 = 1/2.$$

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- Physical meaning: **knowledge of the state of one particle reveals the state of the other.**

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- Uses two conjugate bases  $+ = \{\uparrow, \rightarrow\}$  and  $\times = \{\nearrow, \nwarrow\}$  to establish a secret key between two parties at a distance.
- Suppose they fix the convention that the first vector in each basis represents 0 and the second vector in each basis represents 1.

## A Pictorial Description

Alice's bit	0	1	1	0	1	0	0	1
Alice's basis	+	+	X	+	X	X	X	+
Alice's polarization	↑	→	↙	↑	↙	↗	↗	→
Bob's basis	+	X	X	X	+	X	+	+
Bob's measurement	↑	↗	↙	↗	→	↗	→	→
Public discussion								
Shared Secret key	0		1			0		1



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## Basic Eavesdropping: Measure and Resend

- **No cloning theorem** (Wooters & Zurek, Nature 299, 1982) assures that Eve cannot replicate a particle of unknown state.
- Eve has to measure the photons sent by Alice before sending them on to Bob.
- Since Eve will not know what bases Alice used to encode the bit until after Alice and Bob discuss their measurements, Eve will be forced to guess the basis randomly.

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- If she measures on the incorrect bases, Bob will read a bit incorrectly 50% of the time.
- Given that Eve will choose the measurement basis incorrectly on average 50% of the time, 25% of Bob's measured bits will differ from Alice.
- If Eve has eavesdropped on all the bits, then after  $n$  bit comparisons by Alice and Bob, they will reduce the probability that Eve will go undetected to  $0.75^n$ .



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- If the number of disagreement is more than an acceptable limit then the protocol is aborted.
- Information reconciliation and privacy amplification are performed by Alice and Bob on the remaining  $n$  bits to obtain  $m$  shared key bits.

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- The quantum channel, where any measurement introduces error (**qubits cannot be copied**).  
Note: **because of no-cloning, quantum channel is a perfectly secure communication channel from classical point of view.**
- Orthogonal basis to represent 0 and 1 (**non-orthogonal states cannot be distinguished**).

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- Measurement Device Independent (MDI) Quantum Key Distribution (QKD) idea has been presented very recently by Lo, Curty and Qi (PRL, 2012).
- The idea is to resist detector side channel attacks.
- All the measurements are executed at Eve's end, an untrusted third-party.
- It is natural to consider that Eve herself will try to gather information about the secret key while assisting Alice and Bob.



## Bell States

- The MDI QKD algorithm uses *Bell states*.
- These are two-qubit entangled states.
- Denoted by  $|\Phi^\pm\rangle = \frac{1}{\sqrt{2}}[|00\rangle \pm |11\rangle]$ ,  $|\Psi^\pm\rangle = \frac{1}{\sqrt{2}}[|01\rangle \pm |10\rangle]$ .
- They form an orthogonal basis in a four-dimensional Hilbert space.

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- Alice and Bob announces the bases and discards the mismatched ones.
- For the matched bases, one of Alice and Bob flips the bit, if
  - both are in  $Z$  basis and Eve's outcome are  $|\Psi^\pm\rangle$ , OR
  - both are in  $X$  basis and Eve's outcome are  $|\Phi^-\rangle$  or  $|\Psi^-\rangle$ .

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  - both are in  $X$  basis and Eve's outcome are  $|\Phi^-\rangle$  or  $|\Psi^-\rangle$ .
- Error estimation, information reconciliation and privacy amplification are performed by Alice and Bob on the bits at their ends to obtain the final shared key bits.

## Different Cases

Qubits sent by		Probability (Eve's end)				Flip
Alice	Bob	$ \Phi^+\rangle$	$ \Phi^-\rangle$	$ \Psi^+\rangle$	$ \Psi^-\rangle$	
$ 0\rangle$	$ 0\rangle$	$\frac{1}{2}$	$\frac{1}{2}$	0	0	No
$ 0\rangle$	$ 1\rangle$	0	0	$\frac{1}{2}$	$\frac{1}{2}$	Yes
$ 1\rangle$	$ 0\rangle$	0	0	$\frac{1}{2}$	$\frac{1}{2}$	Yes
$ 1\rangle$	$ 1\rangle$	$\frac{1}{2}$	$\frac{1}{2}$	0	0	No
$ +\rangle$	$ +\rangle$	$\frac{1}{2}$	0	$\frac{1}{2}$	0	No
$ +\rangle$	$ -\rangle$	0	$\frac{1}{2}$	0	$\frac{1}{2}$	Yes
$ -\rangle$	$ +\rangle$	0	$\frac{1}{2}$	0	$\frac{1}{2}$	Yes
$ -\rangle$	$ -\rangle$	$\frac{1}{2}$	0	$\frac{1}{2}$	0	No



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## Eavesdropping: Different Models

- Will Eve work on **each individual qubit** or a set of qubits together?
  - **the first one is called the *incoherent attack*,**
  - **the second one is known as *coherent attack*.**
- Will there be **equal error probability** at Bob's end corresponding to different bases?
  - **If this is indeed equal, then we call it *symmetric*.**
  - **Otherwise, we call it *asymmetric*.**

## How does Eve interact?

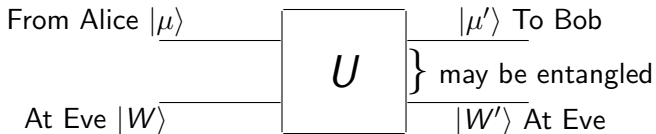


Figure: The model of Eavesdropping

## The unitary interactions at Eve's end

$$\begin{aligned}U|0\rangle|W\rangle &= \sqrt{1-D}|0\rangle|E_{00}\rangle + \sqrt{D}|1\rangle|E_{01}\rangle, \\U|1\rangle|W\rangle &= \sqrt{1-D}|1\rangle|E_{11}\rangle + \sqrt{D}|0\rangle|E_{10}\rangle,\end{aligned}$$

where  $D$  is the disturbance and  $1 - D$  is the fidelity and  $E_{pq}$  is the state of Eve's ancilla qubits after the interaction.

## Expressions for Eve's post-interaction states

$$\begin{aligned}|E_{00}\rangle &= \sqrt{1-D} \frac{|00\rangle + |11\rangle}{\sqrt{2}} + \sqrt{D} \frac{|00\rangle - |11\rangle}{\sqrt{2}}, \\|E_{01}\rangle &= \sqrt{1-D} \frac{|01\rangle + |10\rangle}{\sqrt{2}} - \sqrt{D} \frac{|01\rangle - |10\rangle}{\sqrt{2}}, \\|E_{10}\rangle &= \sqrt{1-D} \frac{|01\rangle + |10\rangle}{\sqrt{2}} + \sqrt{D} \frac{|01\rangle - |10\rangle}{\sqrt{2}}, \\|E_{11}\rangle &= \sqrt{1-D} \frac{|00\rangle + |11\rangle}{\sqrt{2}} - \sqrt{D} \frac{|00\rangle - |11\rangle}{\sqrt{2}}.\end{aligned}$$

It can be shown that Eve can correctly guess the qubit sent by Alice and received by Bob with probability  $\frac{1}{2} + \sqrt{D(1-D)}$ .

## The unitary interactions at Eve's end

$$\begin{aligned}U|0\rangle_A|W\rangle_A &= \sqrt{1-D}|0\rangle_A|E_{00}\rangle_A + \sqrt{D}|1\rangle_A|E_{01}\rangle_A, \\U|1\rangle_A|W\rangle_A &= \sqrt{1-D}|1\rangle_A|E_{11}\rangle_A + \sqrt{D}|0\rangle_A|E_{10}\rangle_A,\end{aligned}$$

where  $D$  is the disturbance and  $1 - D$  is the fidelity and  $E_{pq}$  is the state of Eve's ancilla qubits after the interaction.

## The overall state at Eve's end

An example case: both Alice and Bob send 0.

$$\begin{aligned} & (\sqrt{1-D}|0\rangle_A |E_{00}\rangle_A + \sqrt{D}|1\rangle_A |E_{01}\rangle_A) |0\rangle_B \\ = & \sqrt{\frac{1-D}{2}} |E_{00}\rangle_A |\phi^+\rangle_{AB} + \sqrt{\frac{1-D}{2}} |E_{00}\rangle_A |\phi^-\rangle_{AB} \\ & + \sqrt{\frac{D}{2}} |E_{01}\rangle_A |\psi^+\rangle_{AB} - \sqrt{\frac{D}{2}} |E_{01}\rangle_A |\psi^-\rangle_{AB}. \end{aligned}$$

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## The overall state at Eve's end

An example case: when Alice sends 0 and Bob sends 1.

$$\begin{aligned} & \left( \frac{1-D}{\sqrt{2}} |F_{0011}\rangle + \frac{D}{\sqrt{2}} |F_{0110}\rangle \right) \Psi_{AB}^+ + \left( \frac{1-D}{\sqrt{2}} |F_{0011}\rangle - \frac{D}{\sqrt{2}} |F_{0110}\rangle \right) \\ & + \left( \sqrt{\frac{D(1-D)}{2}} |F_{0111}\rangle + \sqrt{\frac{D(1-D)}{2}} |F_{0010}\rangle \right) \Phi_{AB}^+ \\ & + \left( \sqrt{\frac{D(1-D)}{2}} |F_{0010}\rangle - \sqrt{\frac{D(1-D)}{2}} |F_{0111}\rangle \right) \Phi_{AB}^- \end{aligned}$$

Here,  $|F_{pqrs}\rangle = |E_{pq}\rangle_A |E_{rs}\rangle_B$ .

## The effective disturbance at Alice and Bob's end

### Proposition

*When Eve eavesdrops on both the sides, the effective disturbance at Alice and Bob's end is given by  $\Delta(D) = 2D(1 - D)$ .*

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## Eve's Likelihoods $P(V = v \mid A = a, B = b)$

$V$	$A = 0,$ $B = 0$	$A = 0,$ $B = 1$	$A = 1,$ $B = 0$	$A = 1,$ $B = 1$
00	$(1 - D)d$	$(1 - D)d$	$(1 - D)d'$	$(1 - D)d'$
01	$Dd'$	$Dd'$	$Dd$	$Dd$
10	$Dd$	$Dd$	$Dd'$	$Dd'$
11	$(1 - D)d'$	$(1 - D)d'$	$(1 - D)d$	$(1 - D)d$

Here  $d = \frac{1}{2} + \sqrt{D(1 - D)}$  and  $d' = \frac{1}{2} - \sqrt{D(1 - D)}$ .

## Success Probability Expressions

### Theorem

*The optimal success probability of Eve in guessing the bit sent by Alice is given by  $\frac{1}{2} + \sqrt{D(1-D)}$ .*

### Corollary

*Success probability for guessing both Alice's and Bob's bit is  $P_1(D) = \frac{1}{2}d = \frac{1}{4} + \frac{1}{2}\sqrt{D(1-D)}$ .*

# Eve's Likelihoods $P(V = v \mid A = a, B = b)$

$V$	$A = 0, B = 0$	$A = 0, B = 1$	$A = 1, B = 0$	$A = 1, B = 1$
0000	$\frac{1}{4}(1 - D)^2 d_+$	$\frac{1}{4}(1 - D)^2 d_2$	$\frac{1}{4}(1 - D)^2 d_2$	$\frac{1}{4}(1 - D)^2 d_-$
0001	$\frac{1}{4}D(1 - D)d_2$	$\frac{1}{4}D(1 - D)d_+$	$\frac{1}{4}D(1 - D)d_-$	$\frac{1}{4}D(1 - D)d_2$
0010	$\frac{1}{4}D(1 - D)d_+$	$\frac{1}{4}D(1 - D)d_2$	$\frac{1}{4}D(1 - D)d_2$	$\frac{1}{4}D(1 - D)d_-$
0011	$\frac{1}{4}(1 - D)^2 d_2$	$\frac{1}{4}(1 - D)^2 d_+$	$\frac{1}{4}(1 - D)^2 d_-$	$\frac{1}{4}(1 - D)^2 d_2$
...	...	...	...	...
1100	$\frac{1}{4}(1 - D)^2 d_2$	$\frac{1}{4}(1 - D)^2 d_-$	$\frac{1}{4}(1 - D)^2 d_+$	$\frac{1}{4}(1 - D)^2 d_2$
1101	$\frac{1}{4}D(1 - D)d_-$	$\frac{1}{4}D(1 - D)d_2$	$\frac{1}{4}D(1 - D)d_2$	$\frac{1}{4}D(1 - D)d_+$
1110	$\frac{1}{4}D(1 - D)d_2$	$\frac{1}{4}D(1 - D)d_-$	$\frac{1}{4}D(1 - D)d_+$	$\frac{1}{4}D(1 - D)d_2$
1111	$\frac{1}{4}(1 - D)^2 d_-$	$\frac{1}{4}(1 - D)^2 d_2$	$\frac{1}{4}(1 - D)^2 d_2$	$\frac{1}{4}(1 - D)^2 d_+$

Here  $d_{\pm} = (\sqrt{1 - D} \pm \sqrt{D})^4$  and  $d_2 = (1 - 2D)^2$ .

## Success Probability Expressions

### Theorem

*The optimal success probability of Eve in guessing a pair of bits sent by Alice and Bob is given by*

$$P_2(D) = \frac{1}{4} + D(1 - D) + \sqrt{D(1 - D)}.$$

### Corollary

*Introducing a disturbance  $\Delta$ , the optimal success probability of Eve in guessing a pair of bits sent by Alice and Bob is given by*

$$\frac{1}{4} + \frac{\Delta}{2} + \sqrt{\frac{\Delta}{2}}.$$

# Eve's optimal guesses corresponding to her measurement outcomes

$V$	0000	0001	0010	0011	0100	0101	0110	0111
$G_A, G_B$	0, 0	0, 1	0, 0	0, 1	1, 0	1, 1	1, 0	1, 1
$V$	1000	1001	1010	1011	1100	1101	1110	1111
$G_A, G_B$	0, 0	0, 1	0, 0	0, 1	1, 0	1, 1	1, 0	1, 1



# Roadmap

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- 2 Key Distribution using BB84
  - The Protocol
  - Security Issues
- 3 MDI QKD
- 4 Attack Model
  - For BB84
  - For MDI QKD on one side
  - For MDI QKD on both sides
- 5 Eve's Success Probabilities
  - One-sided Eavesdropping
  - Two-sided Eavesdropping
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## Error Location in BB84

- If Alice sends  $|0\rangle$ , then Eve will obtain either  $|E_{00}\rangle$  or  $|E_{01}\rangle$ .
- Eve now measures in  $\{|00\rangle, |01\rangle, |10\rangle, |11\rangle\}$  basis.
- If Eve observes  $|00\rangle$  or  $|11\rangle$ , then she knows that Bob obtained  $|0\rangle$ , i.e., no error has been introduced.
- If Eve observes  $|01\rangle$  or  $|10\rangle$ , then she knows that Bob obtained  $|1\rangle$ , i.e., error has been introduced.
- Thus, in BB84 protocol, Eve can decide with certainty whether her interaction has introduced an error or not.

## Error Location: One-sided Eavesdropping

- In this case Alice and Bob produce the bits independently.
- So, looking at Alice's bit, it is not possible to know what happens in case of Bob's bit.
- Thus, Alice has no advantage.

## Error Location: Two-sided Eavesdropping

A	B	$G_A$	$G_B$	$P(A, B, G_A, G_B)$	Error Guessed by Eve correctly
0	0	0	0	$\frac{1}{16} d_+$	Y
		0	1	$\frac{1}{16} d_2$	N
		1	0	$\frac{1}{16} d_2$	N
		1	1	$\frac{1}{16} d_-$	Y
0	1	0	0	$\frac{1}{16} d_2$	N
		0	1	$\frac{1}{16} d_+$	Y
		1	0	$\frac{1}{16} d_-$	Y
		1	1	$\frac{1}{16} d_2$	N

Similarly,  $AB = 10$  and  $11$  can also be analyzed.

## Probability of Eve's Guessing the Error Location

### Theorem

*Eve can guess whether an error has been introduced between Alice and Bob or not with probability  $\frac{1}{2} + 2D(1 - D)$ .*

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- The attack on BB84 is sharper than that on MDI QKD. But the latter case is more challenging for Eve.
- In terms of location of errors, MDI-QKD leaks less information than BB84.

## Future Work

- To analyze different “device independent” protocols under similar attack models.
- To investigate countermeasures against this type of attacks.