Lattices from Totally Real Number Fields with Large Regulator

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Outline

- Introduction
  - Coding Strategy for the Wiretap Rayleigh Fading Channel
  - Code Design Criterion
- Ideal Lattices
- Some Number fields with Prescribed Ramification
  - Norms and Ramification
- Units and Regulator
- Conclusion

Introduction

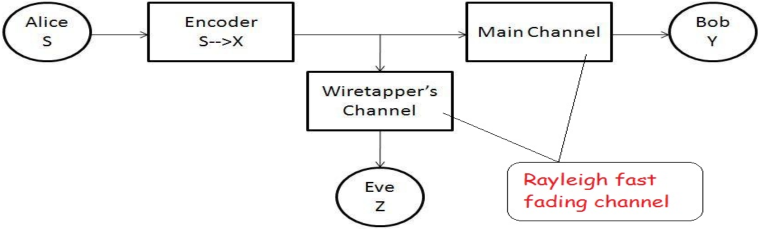

Figure: Wiretap Channel

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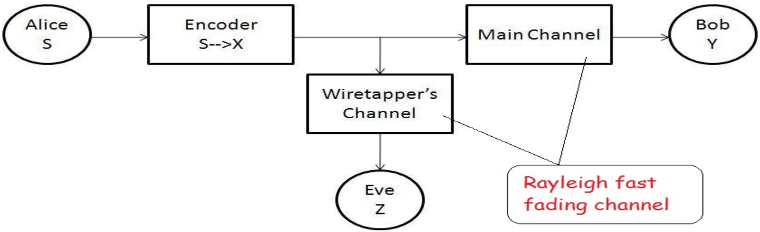

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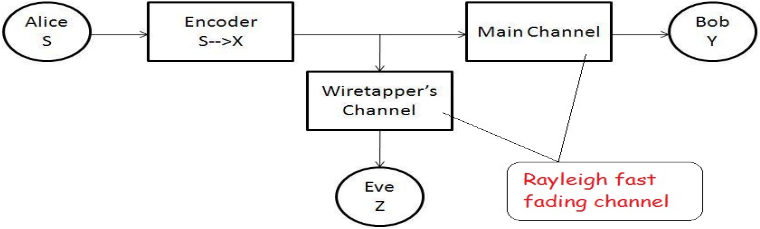

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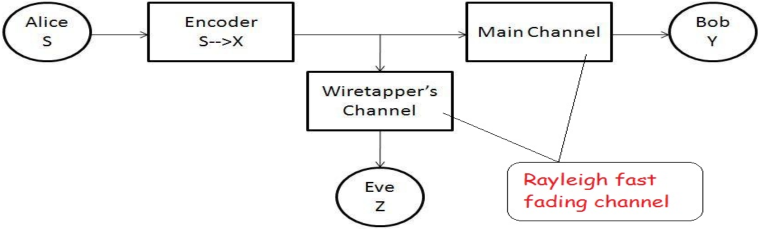

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Coding Strategy for the Wiretap Rayleigh Fading Channel

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Coset encoding:a sublattice Λ_e of Λ_b and partition Λ_b into a union of disjoint cosets of the form

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where $\mathbf{c} = (c_1, c_2, ..., c_n) \in \Lambda_b \subset \mathbb{R}^n$.

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$$\mathbf{x} = \mathbf{r} + \mathbf{c}_{(\mathbf{s})} \in \Lambda_e + \mathbf{c}_{(\mathbf{s})} \Leftrightarrow \text{random vector } \mathbf{r} \in \Lambda_e$$

Code Design Criterion

(J.C. Belfiore and F.Oggier, 2011)

$$\bar{P}_{c,e} \approx (\frac{\gamma_e}{4})^{\frac{n}{2}} \text{Vol}(\Lambda_b) \frac{1}{\gamma_e^{\frac{3}{2}d_{\mathbf{x}}}} \sum_{\mathbf{x}\in\Lambda_e, \mathbf{x}\neq 0} \prod_{x_i \neq 0} \frac{1}{|x_i|^3}$$

where

Λ_b (resp.Λ_e) is the lattice intended for Bob (resp.Eve),
γ_e is Eve's average Signal to Noise Ratio(SNR),
$\mathbf{x} = (x_1, x_2, ..., x_n)$,
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Let K be a totally real number field of degree n, with ring of integers $\mathcal{O}_K$, and real embeddings $\sigma_1, ..., \sigma_n$.

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If $\{\omega_1, \ldots, \omega_n\}$ is a $\mathbb{Z}$-basis of $\mathcal{I}$, the generator matrix M of the corresponding ideal lattice $(\mathcal{I}, q_\alpha) = \{\mathbf{x} = \mathbf{u}M | \mathbf{u} \in \mathbb{Z}^n\}$ is given by

$$M = \begin{pmatrix} \sqrt{\alpha_1}\sigma_1(\omega_1) & \sqrt{\alpha_2}\sigma_2(\omega_1) & \ldots & \sqrt{\alpha_n}\sigma_n(\omega_1) \\ \vdots & \vdots & \ldots & \vdots \\ \sqrt{\alpha_1}\sigma_1(\omega_n) & \sqrt{\alpha_2}\sigma_2(\omega_n) & \ldots & \sqrt{\alpha_n}\sigma_n(\omega_n) \end{pmatrix}$$

where $\alpha_j = \sigma_j(\alpha)$, for all j.

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$\mathbf{x} = (x_1, \ldots, x_n) = (\sqrt{\alpha_1}\sigma_1(x), \ldots, \sqrt{\alpha_n}\sigma_n(x))$ for some $x = \sum_{i=1}^{n} u_i\omega_i \in \mathcal{I} \subseteq \mathcal{O}_K$ where $(u_1, ..., u_n) \in \mathbb{Z}^n$.

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$$= \sum_{x\in\mathcal{I}, x\neq 0} \frac{1}{(N_{K/\mathbb{Q}}(\alpha))^{\frac{3}{2}} |N_{K/\mathbb{Q}}(x)|^3}$$

where $\alpha_j = \sigma_j(\alpha)$, for all j and $N_{K/\mathbb{Q}}(\beta) = \prod_{i=1}^{n} \sigma_i(\beta)$ for $\beta \in K$.

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Thus $x' \in \mathcal{I} = (\beta)\mathcal{O}_K$, $N_{K/\mathbb{Q}}(x') = N_{K/\mathbb{Q}}(\beta)N_{K/\mathbb{Q}}(x)$ for some $x \in \mathcal{O}_K$.

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Hence,

$$\sum_{x' \in \mathcal{I}, x' \neq 0} \frac{1}{|N_{K/\mathbb{Q}}(x')|^3} \Rightarrow \sum_{x \in \mathcal{O}_K, x \neq 0} \frac{1}{|N_{K/\mathbb{Q}}(x)|^3}. \qquad (1)$$

$$
\begin{aligned}
\sum_{x\in\mathcal{O}_K, x\neq 0} \frac{1}{|N_{K/\mathbb{Q}}(x)|^3} &= \frac{A_1}{1^3} + \frac{A_2}{2^3} + \frac{A_{2^2}}{(2^2)^3} + \frac{A_{2^3}}{(2^3)^3} + \ldots \\
&+ \frac{A_3}{3^3} + \frac{A_{3^2}}{(3^2)^3} + \frac{A_{3^3}}{(3^3)^3} + \ldots \\
&+ \frac{A_5}{5^3} + \frac{A_{5^2}}{(5^2)^3} + \frac{A_{5^3}}{(5^3)^3} + \ldots \\
&+ \frac{A_7}{7^3} + \frac{A_{7^2}}{(7^2)^3} + \frac{A_{7^3}}{(7^3)^3} + \ldots
\end{aligned}
$$

where A_i refers to number of algebraic integers with a norm of $\pm i$.

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Instead we will consider in analysing the following

$$\sum_{x \in \mathcal{O}_K \cap \mathcal{R}, x \neq 0} \frac{1}{|N_{K/\mathbb{Q}}(x)|^3}$$

where $\mathcal{R}$ decides the shape of the finite constellation.

Norms and Ramification

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Integers with norms at least 2 depend on

- ramification in K
- the class number of K
- density of units

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$$N(p\mathcal{O}_K) = N(\prod_{i=1}^{g} \mathfrak{p}_i^{e_i}) = \prod_{i=1}^{g} N(\mathfrak{p}_i)^{e_i} = |N_{K/\mathbb{Q}}(p)| = p^n$$

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We can further identify a generator with norm p.

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Example

If 2 is inert prime, then $x \in \mathcal{O}_K$ with $N(x) = 2^k$ for $k \geq n$.

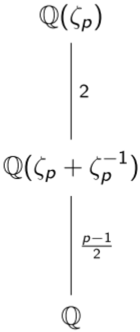

Figure: Cyclotomic Field and its Maximal Real Subfield

Theorem

(D.A.Marcus,1977)
Let q be a rational prime different from p, then q is unramified in $\mathbb{Q}(\zeta_p)$ and in fact

$$(q)\mathbb{Z}[\zeta_p] = \mathfrak{q}_1 ... \mathfrak{q}_g$$

with mutually distinct prime ideals $\mathfrak{q}_i$ and each of inertial degree $f = f(\mathfrak{q}_i/q)$ equal to the order of q in $(\mathbb{Z}/p)^\times$,i.e., f is the least natural number such that

$$q^f \equiv 1 \pmod{p}.$$

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Lemma

Suppose that $p = 2p' + 1$, where both p and p' are prime (such a prime p' is called a Sophie Germain prime). Then the primes smaller than p are inert in $\mathbb{Q}(\zeta_p + \zeta_p^{-1})$.

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Example

Consider $\mathbb{Q}(\zeta_{23})$, with $23 = 2 \cdot 11 + 1$. The primes 2, 3, 5, 7, 11, 13, 17, 19 are all inert in $\mathbb{Q}(\zeta_{23} + \zeta_{23}^{-1})$.

Units and Regulator

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Definition

Given a basis $e_1, \ldots, e_r$ for the group of units modulo the group of roots of unity. The *regulator* of K is

$$R = |\det(\log|\sigma_i(e_j)|)_{1 \leq i,j \leq r}|,$$

where $|\sigma_i(e_j)|$ denotes the absolute value for the real embeddings, and the square of the complex absolute value for the complex ones.

Theorem

(G.R.Everest, J.H.Loxton, 1993)

Let w *be the number of roots of unity in* L. *The number of units* $U(q)$ *such that* $\max_{1\le i\le d}|\sigma_i(u)| < q$ *in* K *is given by*

$$U(q) = \frac{w(r+1)^r}{Rr!}(\log q)^r + O((\log q)^{r-1-(cR^{2/r})^{-1}})$$

as $q \to \infty$ *and* $c = 6\cdot 2\times 10^{12} d^{10}(1+2\log d)$.

Table: Some totally real number fields K of Cyclotomic Fields.

$K \subset \mathbb{Q}(\zeta_p)$	R	$p(X)$	primes
$\mathbb{Q}(\zeta_{11})$	1.63	$x^5 + x^4 - 4x^3 - 3x^2 + 3x + 1$	11 ramifies
$\mathbb{Q}(\zeta_{31})$	30.36	$x^5 - 9x^4 + 20x^3 - 5x^2 - 11x - 1$	5 splits
$\mathbb{Q}(\zeta_{41})$	123.32	$x^5 - x^4 - 16x^3 - 5x^2 + 21x + 9$	3 splits
$\mathbb{Q}(\zeta_{23})$	1014.31	$x^{11} + x^{10} - 10x^9 - 9x^8 + 36x^7 + 28x^6$	
		$-56x^5 - 35x^4 + 35x^3 + 15x^2 - 6x - 1$	23 ramifies
$\mathbb{Q}(\zeta_{67})$	330512.24	$x^{11} - x^{10} - 30x^9 + 63x^8 + 220x^7 - 698x^6$	
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For the case of degree 5,

$$\frac{2 \cdot 5^4}{4!R}(\log q)^4 = \frac{625}{12R}(\log q)^4$$

yielding respectively

$$\sim 32(\log q)^4, \quad \sim 0.4(\log q)^4$$

for the smallest and biggest regulators shown in Table 1.

Conclusion

- Code design criterion for fast fading channel is analysed in designing the lattice code that provides confusion to the eavedropper.

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- Code design criterion for fast fading channel is analysed in designing the lattice code that provides confusion to the eavedropper.
- Identifying totally real number fields with prescribed ramification and regulator provide some thought in the design of wiretap codes for fast fading channels.

$\sim$ Thank you for your attention! $\sim$