

REPORTS IN INFORMATICS

ISSN 0333-3590

Frequency permutation arrays
within distance one

Torleiv Kløve

REPORT NO 382

January 2009



Department of Informatics
UNIVERSITY OF BERGEN
Bergen, Norway

This report has URL <http://www.ii.uib.no/publikasjoner/texrap/ps/2009-382.ps>

Reports in Informatics from Department of Informatics, University of Bergen, Norway, is available at <http://www.ii.uib.no/publikasjoner/texrap/>.

Requests for paper copies of this report can be sent to:

Department of Informatics, University of Bergen, Høyteknologisenteret,
P.O. Box 7800, N-5020 Bergen, Norway

Abstract

The following problem is considered: how many permutations p of the sequence $\iota = (1, 1, \dots, 1, 2, 2, \dots, 2, \dots, n, n, \dots, n)$, where each element occurs l times, satisfy $|p_i - \iota_i| \leq 1$ for all $i = 1, 2, \dots, ln$? It is shown that this number, denoted by $V_l(n)$, satisfies a linear recurrence of length $l + 1$. A table of values of $V_l(n)$ and generating functions for $l \leq 10$ are given.

1 Introduction

Let $S_{l,n}$ denote the set of permutations of the sequence

$$\iota = (1, 1, \dots, 1, 2, 2, \dots, 2, \dots, n, n, \dots, n),$$

where each element occurs l times, or equivalently, $S_{l,n}$ is the $([l, l, \dots, l], 1)$ constant composition code over $\{1, 2, \dots, n\}$. A *frequency permutation array* with minimum distance d is a subset of $S_{l,n}$ where the distance between the permutations are at least d , or equivalently, a $([l, l, \dots, l], d)$ constant composition code. The distance we consider here is the distance obtained from the l_∞ norm, that is, we require that

$$|p_i - \iota_i| \leq d \text{ for all } i, 1 \leq I \leq ln. \quad (1)$$

The number of permutations in $S_{l,n}$ satisfying (1) we denote by $V_l(d, n)$. Such frequency permutation arrays were studied by Shieh and Tsai [3]. For $l = 1$ we get a *permutation array*, and $V_1(d, n)$ has been studied by a number of authors since the 1950s. Kløve [1] gave a survey of known result as well as many new results. It is well known that $V_1(d, n)$ can be given as the value of the permanent of some matrix, and Shieh and Tsai [3] showed that the same is true for $V_l(d, n)$. The permanent of an $n \times n$ matrix $M = (m_{i,j})$ is defined by

$$\text{per } M = \sum_{p \in S_{1,n}} m_{1,p_1} \cdots m_{n,p_n}. \quad (2)$$

It is known that $V_1(d, n)$ satisfies a linear recurrence in n and that the generating function $\sum_{n=0}^{\infty} V_1(d, n)x^n$ is a rational function.

Shieh and Tsai [3] showed that if M is the $ln \times ln$ matrix defined by

$$m_{i,j} = 1 \text{ if } \left| \left\lceil \frac{i}{l} \right\rceil - \left\lceil \frac{j}{l} \right\rceil \right| \leq d \text{ and } m_{i,j} = 0 \text{ otherwise,}$$

then $V_l(d, n) = \text{per } M / (l!)^n$. The known methods can be used to show that $V_l(d, n)$ satisfies some linear recurrence. However, the actual determination of the recurrence will probably be quite complicated in general. In this paper we will study the recursion for $d = 1$ (and all l). For convenience, we drop d from the notation and write $V_l(n) = V_l(1, n)$.

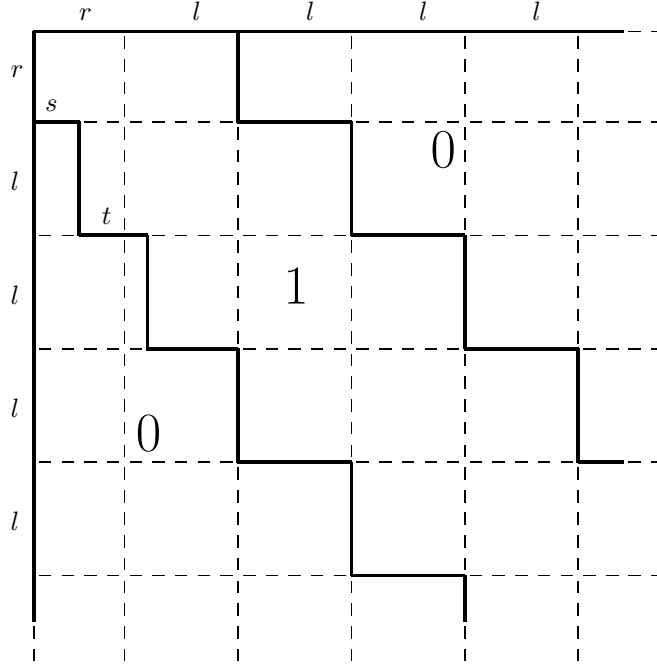
2 The main result

Suppose that l is given. We will define a class of matrices $M_l(n, r, s, t)$, given by four parameters, n, r, s, t , where

$$n \geq 1, \quad 1 \leq r \leq l, \quad 0 \leq s \leq r, \quad 0 \leq t \leq l, \text{ and } r \leq s + t \leq r + l.$$

The matrix M given by Shieh and Tsai [3] mentioned above is $M_l(n - 1, l, 0, l)$ in this notation. The matrix $M_l(n, r, s, t)$ is an $(ln + r) \times (ln + r)$ matrix obtained by deleting the first $l - r$ rows and columns of $M_l(n, l, 0, l)$ and in addition some of

Figure 1: The form of the matrix $M_l(n, r, s, t)$. The upper left corner is shown.



the remaining elements are changed from 1 to 0. Explicitly, $M_l(n, r, s, t) = (m_{i,j})$ where $1 \leq i, j \leq l(n-1) + r$, and where $m_{i,j} = 1$ if (and only if)

$$\begin{aligned} &1 \leq i \leq r \text{ and } 1 \leq j \leq l + r, \\ &r + 1 \leq i \leq l + r \text{ and } s + 1 \leq j \leq 2l + r, \\ &l + r + 1 \leq i \leq 2l + r \text{ and } s + t + 1 \leq j \leq 3l + r, \\ &(a + 1)l + r + 1 \leq i \leq (a + 2)l + r \text{ and } al + 1 \leq j \leq (a + 2)l + r \end{aligned}$$

for $2 \leq a \leq n - 2$,

$$(n - 1)l + r + 1 \leq i \leq nl + r \text{ and } (n - 1)l + 1 \leq j \leq nl + r.$$

The form of $M(n, r, s, t)$ is illustrated in Fig. 1. The 1s are located symmetrically around the main diagonal, except for the first $r + l$ rows and columns.

Let $A_l(n, r, s, t) = \text{per } M_l(n, r, s, t)$. Our main goal is then to determine

$$V_l(n) = A_l(n - 1, l, 0, l) / (l!)^n.$$

We will show that $V_l(n)$ satisfies a linear recursion in n of degree l . We also find its generating function.

We see that the first row of $M_l(n, r, s, t)$ starts with $r + l$ ones (and the remaining elements in the first row are zero). Expanding the permanent by the first row, we therefore get an expression for $A_l(n, r, s, t)$ in terms of some $A_l(n, r - 1, s', t')$ if $r > 1$ and in terms of some $A_l(n - 1, l, t', l - t')$ if $r = 1$. To simplify the notation a little bit, we let

$$B_l(n, u) = A_l(n - 1, l, u, l - u),$$

where we for now assume that n is fixed and we will show that all $A_l(n, r, s, t)$ can be expressed in terms of the $B_l(n, u)$, $u = 0, 1, \dots, l$. In this notation, $V_l(n) = B_l(n, 0)/(l!)^n$.

First consider $r = 1$. Then $0 \leq s \leq 1$. If $s = 1$, the first column contains only one 1 and we get

$$A_l(n, 1, 1, t) = B_l(n, t). \quad (3)$$

If $s = 0$, the first t columns are identical and the next $l + 1 - t$ are also identical. Hence we get

$$A_l(n, 1, 0, t) = t B_l(n, t - 1) + (l + 1 - t) B_l(n, t). \quad (4)$$

Now consider $r > 1$. Of the $r + l$ first columns, the first s are identical, the next t are identical, and the final $r + l - s - t$ are identical. Hence, expanding by the first row, we get

$$\begin{aligned} A_l(n, r, s, t) &= s A_l(n, r - 1, s - 1, t) + t A_l(n, r - 1, s, t - 1) \\ &\quad + (r + l - s - t) A_l(n, r - 1, s, t). \end{aligned} \quad (5)$$

Using (3), (4), and (5) repeatedly, we can determine all $A_l(n, r, s, t)$. We can in fact find an explicit expression for $A_l(n, r, s, t)$ in general.

Lemma 1. *For all r, s, t we have*

$$A_l(n, r, s, t) = r! \sum_{u=0}^t \binom{t}{u} \binom{r+l-s-t}{l-u} B_l(n, u). \quad (6)$$

Proof. To show the lemma, we first show that (3) and (4) are satisfied. For $r > 1$, we use (5) and induction.

First, (3). For $r = s = 1$, the sum in (6) is

$$1! \sum_{u=0}^t \binom{t}{u} \binom{l-t}{l-u} B_l(n, u) = B_l(n, t)$$

since all the other terms are zero. By (3), this is $A_l(n, 1, 1, t)$.

For $r = 1$ and $s = 0$, there are two nonzero terms in the sum, namely for $u = t$ and $u = t - 1$, and so

$$\begin{aligned} &1! \sum_{u=0}^t \binom{t}{u} \binom{l+1-t}{l-u} B_l(n, u) \\ &= t \cdot 1 \cdot B_l(n, t - 1) + 1 \cdot (l + 1 - t) \cdot B_l(n, t) \\ &= A_l(n, 1, 0, t) \end{aligned}$$

by (4).

For $r > 1$, we assume that the expression (6) is true for $r - 1$. Then we get

$$\begin{aligned}
A_l(n, r, s, t) &= s A_l(n, r - 1, s - 1, t) + t A_l(n, r - 1, s, t - 1) \\
&\quad + (r + l - r - t) A_l(n, r - 1, s, t) \\
&= s(r - 1)! \sum_{u=0}^t \binom{t}{u} \binom{r + l - s - t}{l - u} B_l(n, u) \\
&\quad + t(r - 1)! \sum_{u=0}^{t-1} \binom{r + l - s - t}{l - u} B_l(n, u) \\
&\quad + (r + l - s - t)(r - 1)! \sum_{u=0}^t \binom{t}{u} \binom{r - 1 + l - s - t}{l - u} B_l(n, u) \\
&= (r - 1)! \sum_{u=0}^t \binom{t}{u} \binom{r + l - s - t}{l - u} B_l(n, u) \\
&\quad \cdot \left\{ s + (t - u) + (r - t - s + u) \right\} \\
&= r! \sum_{u=0}^t \binom{t}{u} \binom{r + l - s - t}{l - u} B_l(n, u).
\end{aligned}$$

This completes the induction and the proof of the lemma. \square

Choosing $r = l$ and $t = l - s$ in (6) and reducing n by one, we get the following corollary.

Corollary 1.

$$B_l(n, s) = l! \sum_{u=0}^{l-s} \binom{l-s}{u} \binom{l}{u} B_l(n-1, u). \quad (7)$$

Since

$$B_l(0, 0) = 1 \text{ and } B_l(0, s) = 0 \text{ for } s = 1, 2, \dots, l,$$

(7) and induction proves the following.

Corollary 2. For all l, s and n , $(l!)^n$ divides $B_l(n, s)$.

Let

$$b_l(n, s) = \frac{B_l(n, s)}{(l!)^n}.$$

By Corollary 2, the $b_l(n, s)$ are integers, and by Corollary 1 they satisfy the following recursion:

$$b_l(n, s) = \sum_{u=0}^{l-s} \binom{l-s}{u} \binom{l}{u} b_l(n-1, u). \quad (8)$$

We note that $V_l(n) = b_l(n, 0)$.

Let $T_l = (t_{u,s})$ denote the $(l+1) \times (l+1)$ matrix defined by

$$t_{u,s} = \binom{l-s}{u} \binom{l}{u}.$$

Then

$$(b_l(n, 0), b_l(n, 1), \dots, b_l(n, l)) = (b_l(n-1, 0), b_l(n-1, 1), \dots, b_l(n-1, l)) T_l.$$

In particular, if

$$f(x) = \sum_{i=0}^l c_i x^i \quad (9)$$

is the characteristic polynomial of T_l , then

$$\sum_{i=0}^l c_i b_l(n+i, s) = 0.$$

Hence we have proven the following theorem.

Theorem 1. *For each $l \geq 1$, the sequence $V_l(0), V_l(1), V_l(2), \dots$ satisfies the linear recurrence of length $l+1$ determined by the characteristic polynomial of the matrix*

$$T_l = \left[\binom{l-s}{u} \binom{l}{l-u} \right]_{0 \leq u \leq l, 0 \leq s \leq l}.$$

The same is true for the sequences $b_l(0, s), b_l(1, s), b_l(2, s), \dots$ for $0 \leq s \leq l$.

Computing $b_l(n, s)$ for some range of n , an efficient way is to use (8), with the initial values $b_l(0, 0) = 1$ and $b_l(0, s) = 0$ for $1 \leq s \leq l$. The tables of $V_l(n) = b_l(n, 0)$ in the first appendix are computed this way.

Keeping n fixed

In the discussion above we considered the sequence $b_l(n, s)$ for fixed l and s . Alternatively, we can keep n and s fixed. We can get explicit expressions (involving sums).

We have, $b_l(0, 0) = 1$ and $b_l(0, s) = 0$ for $0 \leq s \leq l$. From (8) we get, for $0 \leq s \leq l$,

$$\begin{aligned} b_l(1, s) &= 1, \\ b_l(2, s) &= \sum_{u=0}^{l-s} \binom{l-s}{u} \binom{l}{l-u} = \binom{2l-s}{l}, \\ b_l(3, s) &= \sum_{u=0}^{l-s} \binom{l-s}{u} \binom{l}{l-u} \binom{2l-u}{l}, \\ b_l(4, s) &= \sum_{u=0}^{l-s} \binom{l-s}{u} \binom{l}{l-u} \sum_{j=0}^{l-u} \binom{l-u}{j} \binom{l}{l-j} \binom{2l-j}{l} \\ &= \sum_{j=0}^l \binom{l}{l-j} \binom{2l-j}{l} \sum_{u=0}^l \binom{l-s}{u} \binom{l}{l-u} \binom{l-u}{j} \\ &= \sum_{j=0}^l \binom{l}{l-j} \binom{2l-j}{l} \sum_{u=0}^l \binom{l-s}{u} \binom{l}{j} \binom{l-j}{u} \\ &= \sum_{j=0}^l \binom{l}{j}^2 \binom{2l-j}{l} \sum_{u=0}^l \binom{l-s}{u} \binom{l-j}{l-j-u} \\ &= \sum_{j=0}^l \binom{l}{j}^2 \binom{2l-j}{l} \binom{2l-s-j}{l-j}. \end{aligned}$$

We note that the sequence $\{b_l(3, 0)\}$ is sequence A005258 in [4] and $\{b_l(4, 0)\}$ is sequence A005258 in [4].

3 Generating functions

Let $h_{l,s}(x) = \sum_{n=0}^{\infty} b_l(n, s)x^n$ be the generating function of $b_l(n, s)$. Once we have the values of $b_l(n, s)$ for $n = 0, 1, \dots, l$, then by the standard theory of linear recurrences, if $g_l(x) = \sum_{i=0}^l \gamma_{l,i}x^i$, where $\gamma_{l,i} = c_{l-i}$ (here the c_i are the coefficient of the characteristic polynomial (9) of T_l), then

$$f_{l,s}(x) = g_l(x) \sum_{n=0}^{\infty} b_l(n, s)x^n$$

is a polynomial of degree less than l , and so

$$h_{l,s}(x) = \frac{f_{l,s}(x)}{g_l(x)}.$$

Tables of $g_l(x)$ and $f_{l,s}(x)$ for $0 \leq s < l \leq 10$ are given in the second appendix. Note that by (8), $b_l(n, l) = b_l(n-1, 0)$. In particular, this implies that

$$f_{l,l}(x) = xf_{l,0}(x).$$

Therefore, we have not included $f_{l,l}(x)$ in the tables.

We note that $\gamma_{l,l+1} = \det(M) = (-1)^{\lceil \frac{l+1}{2} \rceil} \prod_{i=0}^l \binom{l}{i}$. We use the notation $f_{l,s}(x) = \sum_{i=0}^l \phi_{l,s,i}x^i$. Some observed properties of these polynomials for the range we have computed ($l \leq 10$) are given below. These properties may be true in general, but we have not studied this further.

- For $0 \leq s < l$, we have $\deg f_{l,s}(x) = l-1$, that is, $\phi_{l,s,l} = 0$.
- $\phi_{l,s,l-1} = (-1)^{s+1} \frac{\gamma_{l,l+1}}{\binom{l}{s}}$,
- $(-1)^{\lceil \frac{l}{2} \rceil} \gamma_{l,i} > 0$ for $0 \leq i \leq l+1$,
- $(-1)^{\lceil \frac{l}{2} \rceil} \phi_{l,0,i} > 0$ for $0 \leq i \leq l-1$,
- $(-1)^{\lceil \frac{l-1}{2} \rceil} \phi_{l,1,i} > 0$ for $0 \leq i \leq l-1$.

It is also possible to compute the $h_{l,s}(x)$ more directly. Let $\delta_{s,0}$ denote the Kronecker delta, that is $\delta_{0,0} = 1$ and $\delta_{s,0} = 0$ for $s \neq 0$. From (8) and the fact that $b_l(0, s) = \delta_{s,0}$, we get

$$\begin{aligned} h_{l,s}(x) &= \delta_{s,0} + \sum_{n=1}^{\infty} b_l(n, s)x^n \\ &= \delta_{s,0} + \sum_{n=1}^{\infty} x^n \sum_{u=0}^l \binom{l-s}{u} \binom{l}{u} b_l(n-1, u) \\ &= \delta_{s,0} + \sum_{u=0}^l \binom{l-s}{u} \binom{l}{u} \sum_{n=1}^{\infty} b_l(n-1, u)x^n \\ &= \delta_{s,0} + \sum_{u=0}^l \binom{l-s}{u} \binom{l}{u} x h_{l,u}(x). \end{aligned}$$

We can solve this set of $l+1$ equations and obtain the $h_{l,s}(x)$. In terms of matrices, we can write the equation set

$$(h_{l,0}(x), h_{l,1}(x), \dots, h_{l,l}(x)) (I - xT_l) = (1, 0, \dots, 0),$$

where I the $(l+1) \times (l+1)$ identity matrix. Hence we get the following theorem.

Theorem 2. For $l \geq 1$, $(h_{l,0}(x), h_{l,1}(x), \dots, h_{l,l}(x))$ is the first row of

$$(I - xT_l)^{-1}.$$

References

- [1] T. Kløve, "Spheres of Permutations under the Infinity Norm - Permutations with limited displacement," Department of Informatics, University of Bergen, Report No. 376, November 2008.
- [2] T. Kløve, T.-T. Lin, S.-C. Tsai, W.-G. Tzeng, "Efficient encoding and decoding with permutation arrays under infinity norm." Manuscript 2008.
- [3] M.-Z. Shieh and S.-C. Tsai, "Decoding frequency permutation arrays under infinity norm." Manuscript submitted to ISIT 2009.
- [4] N. J. A. Sloane, "The On-Line Encyclopedia of Integer Sequences," <http://www.research.att.com/njas/sequences/>

Appendix, tables of $V_l(n)$

For each $l \leq 10$, we have included the values of $V_l(n)$ for $n \leq 20$ (provided the values can fit into one line, that is, have at most 70 digits). For $l = 1$ the values are the Fibonacci numbers.

$V_1(n)$:

1, 2,
3, 5,
8, 13,
21, 34,
55, 89,
144, 233,
377, 610,
987, 1597,
2584, 4181,
6765, 10946

$V_2(n)$:

1,
6,
19,
73,
264,
973,
3565,
13086,
48007,
176149,
646296,
2371321,
8700553,
31923030,
117128107,
429752305,
1576795176,
5785386229,
21227039605,
77883687150

$V_3(n)$:

1,
20,
147,
1445,
13040,
120685,
1108677,
10207204,
93913687,
864237977,
7952680800,
73181430953,
673420201433,
6196864537204,
57024000600459,
524739088121629,
4828687880190544,
44433943310167925,
408884434208138877,
3762584822284713668

$V_4(n)$:

1,
70,
1251,
33001,
778840,
18979501,
458283501,
11095029310,
268407261751,
6494628568501,
157140210849480,
3802140006795001,
91995514253325001,
2225900960991467350,
53857332463462696251,
1303118450607052410001,
31529924703022030340920,
762890097390817464052501,
18458696127973957622400501,
446621950121084699577258190

$V_5(n)$:

1,
252,
11253,
819005,
51955008,
3426862513,
223555999521,
14629457476284,
956506304490805,
62554067142502589,
4090653574960950144,
267509059547920101233,
17493706803334667625377,
1143999628072529268292860,
74811735155413338440277237,
4892306161272288597410401597,
319931875193202258763349275584,
20921913394524835983550484261681,
1368186457085199494502511814687265,
89472418108917375447330264236027196

$V_6(n)$:

1,
924,
104959,
21460825,
3725735664,
677246940469,
121533839931421,
21888321066877356,
3938124367151767675,
708743612574056867221,
127542377579526691799136,
22952473176498367236198409,
4130491868245063608463767529,
743318274070910835570229477500,
133766583990791214326410385255719,
24072462709511521506885699191291809,
4332049316570097565180842830546688336,
779590011562528975268061048428295217261,
140294013270250427479554925525141312195525,
25247129752917525999708412524109996498245516

$V_7(n)$:

1,
3432,
1004307,
584307365,
281152383072,
142648550730349,
71317516434003525,
35801871486260585544,
17952292795232366702935,
9004755701212513017959321,
4516329337951820803551129024,
2265217670079812211495904106153,
1136138585458716600268637214790361,
569840751629792139416688379526917800,
285808716115094273711695504392530084619,
143349934164043810258573400532731379769885,
71898446844750547021445142046885306393733664,
36061312115189434126727298190818148390989166709,
18086875091406590626472543912280627941050344054525,
9071634715222858104813326808457630531494707987562120,
4549959900223147403498497880417205059109504135655606127

$V_8(n)$:

1,
12870,
9793891,
16367912425,
22035642730200,
31502711112595885,
44295290280756102637,
62569846044620145720126,
88271313599568437690474935,
124573943375568956577577001653,
175789302808103093628255125104200,
248067238178807267808839236968317689,
350060517174005010391130870435408896201,
493989534503292754868353691591517256079190,
697095240068764884987582473125418987303636923,
983708797327298788273493429586552873326822654353,
1388164629137817856436379753509747271807084096142776,
1958914128119468029901132307723927945190375274909584405,
2764329581571820330816482241775044315481460863833230460405,
3900894854445014677762379303942960864562913439627571254125518

$V_9(n)$:

1,
48620,
96918753,
468690849005,
1778348612268800,
7216179759611167513,
28745299946487169993521,
115091311345187072928558460,
460158123142172721285948269305,
1840518625861555681730223351316589,
7360831026278314234828047495128099520,
29439220260831988191360006001082378535233,
117739509112349463723345100673103757213845377,
470889616655296059827467951144765028870328114700,
1883283660494579840168760179636870546935247688814737,
7532036864824211990864954832287836353353586148141111597,
30123755570935847905978484025312642337692092111082459475840,
120477458857172786944857749445901131162481044090709350499496681,
481839590235214778858627919090156511553286141181313928510263717265,
1927077423029371736699043021760573914460313105578776085339808319254300

$V_{10}(n)$:

1,
184756,
970336269,
13657436403073,
146898506147371264,
1701765313758012953473,
19316611911260692278259065,
220494323800641046774094312836,
2513010804923623461613325281197757,
28653350779401783399803031179454720649,
326667439282263161656690619517700109863296,
3724347552175232603674450917904230607245244321,
42461049822986464700029958979381560005857008301553,
484096904839865210883595839396855576288524521889773780,
5519167755806190714476276496104887012101313653960556653357,
62923803056370379198572727452389636863587280852479303520268305,
717391637051748688222824953419630961585872405673287610725811666176,
8178952050692287035831370410106122047491833022842126660028825762249729

Appendix, generating functions

g_1	$1 - x - x^2$
$f_{1,0}$	1

g_2	$1 - 3x - 3x^2 + 2x^3$
$f_{2,0}$	$1 - 2x$
$f_{2,1}$	x

g_3	$1 - 7x - 22x^2 + 15x^3 + 9x^4$
$f_{3,0}$	$1 - 6x - 9x^2$
$f_{3,1}$	$x + 3x^2$
$f_{3,2}$	$x - 3x^2$

g_4	$1 - 19x - 139x^2 + 314x^3 + 184x^4 - 96x^5$
$f_{4,0}$	$1 - 18x - 88x^2 + 96x^3$
$f_{4,1}$	$x + 16x^2 - 24x^3$
$f_{4,2}$	$x - 4x^2 + 16x^3$
$f_{4,3}$	$x - 14x^2 - 24x^3$

g_5	$1 - 51x - 1026x^2 + 5375x^3 + 10575x^4 - 5250x^5 - 2500x^6$
$f_{5,0}$	$1 - 50x - 825x^2 + 2750x^3 + 2500x^4$
$f_{5,1}$	$x + 75x^2 - 450x^3 - 500x^4$
$f_{5,2}$	$x + 5x^2 + 150x^3 + 250x^4$
$f_{5,3}$	$x - 30x^2 - 25x^3 - 250x^4$
$f_{5,4}$	$x - 45x^2 - 450x^3 + 500x^4$

g_6	$1 - 141x - 7644x^2 + 111179x^3 + 498171x^4 - 842310x^5 - 369900x^6 + 162000x^7$
$f_{6,0}$	$1 - 140x - 6861x^2 + 78210x^3 + 207900x^4 - 162000x^5$
$f_{6,1}$	$x + 321x^2 - 7560x^3 - 30150x^4 + 27000x^5$
$f_{6,2}$	$x + 69x^2 + 798x^3 + 9540x^4 - 10800x^5$
$f_{6,3}$	$x - 57x^2 + 882x^3 - 4320x^4 + 8100x^5$
$f_{6,4}$	$x - 113x^2 - 1974x^3 - 540x^4 - 10800x^5$
$f_{6,5}$	$x - 134x^2 - 4935x^3 + 32850x^4 + 27000x^5$

g_7	$1 - 393x - 59193x^2 + 2322404x^3 + 28870212x^4 - 109325076x^5 - 164634169x^6 + 64790985x^7 + 26471025x^8$
$f_{7,0}$	$1 - 392x - 56154x^2 + 1918742x^3 + 17656954x^4 - 38319960x^5 - 26471025x^6$
$f_{7,1}$	$x + 1323x^2 - 110005x^3 - 1683787x^4 + 4934055x^5 + 3781575x^6$
$f_{7,2}$	$x + 399x^2 - 3381x^3 + 268569x^4 - 1404585x^5 - 1260525x^6$
$f_{7,3}$	$x - 63x^2 + 12789x^3 + 12691x^4 + 641067x^5 + 756315x^6$
$f_{7,4}$	$x - 273x^2 - 5439x^3 - 156751x^4 - 338541x^5 - 756315x^6$
$f_{7,5}$	$x - 357x^2 - 29253x^3 + 307671x^4 - 276115x^5 + 1260525x^6$
$f_{7,6}$	$x - 385x^2 - 46893x^3 + 1148707x^4 + 5870445x^5 - 3781575x^6$

g_8	$1 - 1107x - 466155x^2 + 50943066x^3 + 1698359832x^4$ $-17812454688x^5 - 58970736384x^6 + 80038748160x^7$ $+28677390336x^8 - 11014635520x^9$
$f_{8,0}$	$1 - 1106x - 454392x^2 + 46023712x^3 + 1275963136x^4$ $-9584418816x^5 - 17662754816x^6 + 11014635520x^7$
$f_{8,1}$	$x + 5328x^2 - 1512504x^3 - 79579136x^4$ $+871670016x^5 + 2035740672x^6 - 1376829440x^7$
$f_{8,2}$	$x + 1896x^2 - 189576x^3 + 4129216x^4$ $-150226944x^5 - 518418432x^6 + 393379840x^7$
$f_{8,3}$	$x + 180x^2 + 123456x^3 + 4014976x^4$ $+24159744x^5 + 217061376x^6 - 196689920x^7$
$f_{8,4}$	$x - 612x^2 + 32112x^3 - 4323200x^4$ $+19869696x^5 - 126443520x^6 + 157351936x^7$
$f_{8,5}$	$x - 942x^2 - 157176x^3 + 712768x^4$ $-62243328x^5 + 69543936x^6 - 196689920x^7$
$f_{8,6}$	$x - 1062x^2 - 316056x^3 + 16515520x^4$ $+118315008x^5 + 155947008x^6 + 393379840x^7$
$f_{8,7}$	$x - 1098x^2 - 411768x^3 + 33923680x^4$ $+639330048x^5 - 2611058688x^6 - 1376829440x^7$
g_9	$1 - 3139x - 3733258x^2 + 1141936935x^3 + 105766383267x^4$ $-2954887228278x^5 - 26984014044012x^6 + 79103022450480x^7$ $+98126014508928x^8 - 32338686530304x^9 - 11759522374656x^{10}$
$f_{9,0}$	$1 - 3138x - 3687777x^2 + 1082504250x^3 + 89860199580x^4$ $-2022819810480x^5 - 12869477292672x^6 + 20579164155648x^7$ $+11759522374656x^8$
$f_{9,1}$	$x + 21171x^2 - 19932318x^3 - 3531961260x^4$ $+128735136432x^5 + 1103288522112x^6 - 2141394506496x^7$ $-1306613597184x^8$
$f_{9,2}$	$x + 8301x^2 - 3847878x^3 - 109098090x^4$ $-11363442048x^5 - 188865785808x^6 + 489980098944x^7$ $+326653399296x^8$
$f_{9,3}$	$x + 1866x^2 + 846207x^3 + 196166610x^4$ $-1966952808x^5 + 41823698112x^6 - 184992486336x^7$ $-139994313984x^8$
$f_{9,4}$	$x - 1137x^2 + 798984x^3 - 67272606x^4$ $+3813224040x^5 - 1402059456x^6 + 101107004544x^7$ $+93329542656x^8$
$f_{9,5}$	$x - 2424x^2 - 648891x^3 - 81203310x^4$ $-3419960616x^5 - 23288295744x^6 - 69997156992x^7$ $-93329542656x^8$
$f_{9,6}$	$x - 2919x^2 - 2062908x^3 + 199842228x^4$ $-807352920x^5 + 60553096128x^6 + 34998578496x^7$ $+139994313984x^8$
$f_{9,7}$	$x - 3084x^2 - 3007863x^3 + 585198918x^4$ $+22357365660x^5 - 114671268432x^6 + 190547816256x^7$ $-326653399296x^8$
$f_{9,8}$	$x - 3129x^2 - 3497238x^3 + 903113388x^4$ $+58072548240x^5 - 864203582592x^6 - 2939880593664x^7$ $+1306613597184x^8$

g_{10}	$1 - 8953x - 30255928x^2 + 26167186427x^3$ $+ 6764138451925x^4 - 523127064454800x^5 - 12617955113431500x^6$ $+ 101565785269440000x^7 + 267474303094050000x^8 - 306145784385000000x^9$ $- 93630456360000000x^{10} + 32406091200000000x^{11}$
$f_{10,0}$	$1 - 8952x - 30080125x^2 + 25453146300x^3$ $+ 6170357191500x^4 - 416764490490000x^5 - 7880125909050000x^6$ $+ 43309237545000000x^7 + 61224365160000000x^8 - 32406091200000000x^9$
$f_{10,1}$	$x + 83425x^2 - 255730650x^3$ $- 148914546750x^4 + 17884439536500x^5 + 497095510815000x^6$ $- 3484893302100000x^7 - 5798375604000000x^8 + 3240609120000000x^9$
$f_{10,2}$	$x + 34805x^2 - 62361320x^3$ $- 14123649300x^4 - 336837114000x^5 - 52101666450000x^6$ $+ 573973533000000x^7 + 1200511368000000x^8 - 720135360000000x^9$
$f_{10,3}$	$x + 10495x^2 + 1566720x^3$ $+ 6691267350x^4 - 569339505000x^5 + 1254439980000x^6$ $- 135720663750000x^7 - 408934008000000x^8 + 270050760000000x^9$
$f_{10,4}$	$x - 945x^2 + 9963680x^3$ $+ 316601250x^4 + 344117605500x^5 + 6677031825000x^6$ $+ 30313934100000x^7 + 203364756000000x^8 - 154314720000000x^9$
$f_{10,5}$	$x - 5950x^2 - 144275x^3$ $- 3090409500x^4 - 47843392500x^5 - 8059670100000x^6$ $+ 14074256250000x^7 - 135790830000000x^8 + 128595600000000x^9$
$f_{10,6}$	$x - 7952x^2 - 12506625x^3$ $+ 1230762300x^4 - 321315399000x^5 + 6648804945000x^6$ $- 59447544300000x^7 + 106366932000000x^8 - 154314720000000x^9$
$f_{10,7}$	$x - 8667x^2 - 21619300x^3$ $+ 9550455000x^4 + 599180215500x^5 + 5974944615000x^6$ $+ 149547242250000x^7 - 37614213000000x^8 + 270050760000000x^9$
$f_{10,8}$	$x - 8887x^2 - 26845400x^3$ $+ 17532273900x^4 + 2556502038000x^5 - 74233369890000x^6$ $- 282790683000000x^7 - 559819512000000x^8 - 720135360000000x^9$
$f_{10,9}$	$x - 8942x^2 - 29245875x^3$ $+ 22895839800x^4 + 4681211724000x^5 - 237920095125000x^6$ $- 2909170800900000x^7 + 8460304524000000x^8 + 3240609120000000x^9$